

AN14150

MCX Nx4x Inter-Core Communication

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Application note

Document information

Information	Content
Keywords	Nirvana MCX Nx4x, multi core communication, Mailbox
Abstract	This application note introduces how dual core devices can communicate using the Mailbox interface.



1 Introduction

The MCX Nx4x series microcontrollers combine two Arm Cortex M33 cores with a CoolFlux BSP32, a PowerQuad DSP Co-processor, an NPU, and multiple high-speed connectivity options running at 150 MHz. To support a wide variety of applications, the MCX N-series includes advanced serial peripherals, timers, high-precision analog, and state-of-the-art security features. All MCX Nx4x products include dual-bank flash that supports read while write operation from internal flash. The MCX Nx4x series also supports large external serial memory configurations.

The MCX Nx4x series is a dual-core microcontroller family. CPU0 is the primary Cortex-M33 (ver r0p4-00rel0) processor that supports TrustZone-M, Floating Point Unit (FPU), and Memory Protection Unit (MPU).

The MCX Nx4x device also includes a second instance of Cortex-M33, CPU1 that is the secondary CM33 intended to offload work from the main processor to support special dedicated applications. The configuration of this instance does not include MPU, FPU, DSP, ETM, Trustzone-M, Secure Attribution Unit (SAU) or co-processor interface. SYSTICK is supported on both cores.

Cortex-M33 implements a modified Harvard memory architecture using two 32-bit bus interfaces: the Code and System buses. The bus interfaces are activated by address range and can include both instruction fetches and operand data references on a given bus port. (A traditional Harvard architecture strictly separates instruction fetches and operand data references onto specific bus ports regardless of access address.) The Code bus is typically used for instruction fetching and data accesses of PC-relative data, while the system bus is typically used for operand data references to the on-chip and off-chip memories and peripheral accesses. The bus structure fully supports concurrent instruction fetch and data access, but the Cortex-M33 implementations can generate both types of references on each bus.

2 Dual-core basic mechanism

CPU0 is the main core that boots first and acts as the Master core. CPU1 acts as the Slave core and is held in reset when the device boots up to minimize power consumption. Therefore, to run or debug applications in the CPU1, some code must be executed on CPU0 to initialize CPU1.

In dual-core running mode, CPU0 and CPU1 need to communicate with each other. The MCX Nx 4x provides an Inter-CPU Mailbox module (MAILBOX) to facilitate this communication by using two mechanisms:

- **Interrupts mechanism**

The first mechanism triggers an interrupt, along with a simple message, to another CPU. To generate an interrupt to another CPU, you can set any one of the interrupt flags (up to 32) by using Cortex_M33 (CPU0) Interrupt Set (IRQ0SET) or Cortex_M33 (CPU1) Interrupt Set (IRQ1SET). The other CPU reads Cortex_M33 (CPU0) Interrupt Set (IRQ0SET) and Cortex_M33 (CPU1) Interrupt Set (IRQ1SET) to determine the action it must perform. If more than one IRQn field is 1, the CPU receiving the interrupts reads all of them. After completing the requested service or operation, the CPU that receives the interrupt requests, writes 1 to one or more fields in Cortex_M33 (CPU0) Interrupt Clear (IRQ0CLR) or Cortex_M33 (CPU1) Interrupt Clear (IRQ1CLR) to clear them.

To check whether the operation is complete, the CPU sending the request can check IRQn, or the CPU receiving the original interrupt can use the same mechanism on the reverse to notify the requesting CPU. The user determines the particular meaning of each bit in Cortex_M33 (CPU0) Interrupt (IRQ0) and Cortex_M33 (CPU1) Interrupt (IRQ1) for their specific application, and provides additional information to each request bit. For example, both the CPUs can use predefined memory locations to hold detailed information about the interrupt requests.

- **Mutex mechanism**

The second mechanism is based on a single-bit resource allocation request. The MUTEX[EX] bit defines a mutual exclusion request, and reading that bit provides the status of the resource. If the resource is available, it can be reserved for the process that reads the register. The user defines the resource control, and after

getting access to the resource, can write 1 to MUTEX[EX] to signal to the other processor that the shared resource is now available. For more information, see the MUTEX[EX] description in the MCX Nx 4x Reference Manual.

In summary, the Inter-CPU Mailbox provides the below features:

- Provides interprocessor communication, allowing multiple CPUs to share resources and communicate with each other in a simple manner.
- Enables each CPU to generate up to 32 user-defined interrupts.
- Enables each CPU to claim a shared resource using the MUTEX register bit.

3 Inter-CPU Mailbox

This Inter-CPU Mailbox is an AHB slave peripheral with a single clock domain and reset. The Inter-CPU Mailbox has one mutex register, and three 32-bit registers per CPU (IRQ0/1, IRQ0/1SET, iIRQ0/1CLR).

Figure 1 shows the block diagram of the Inter-CPU Mailbox:

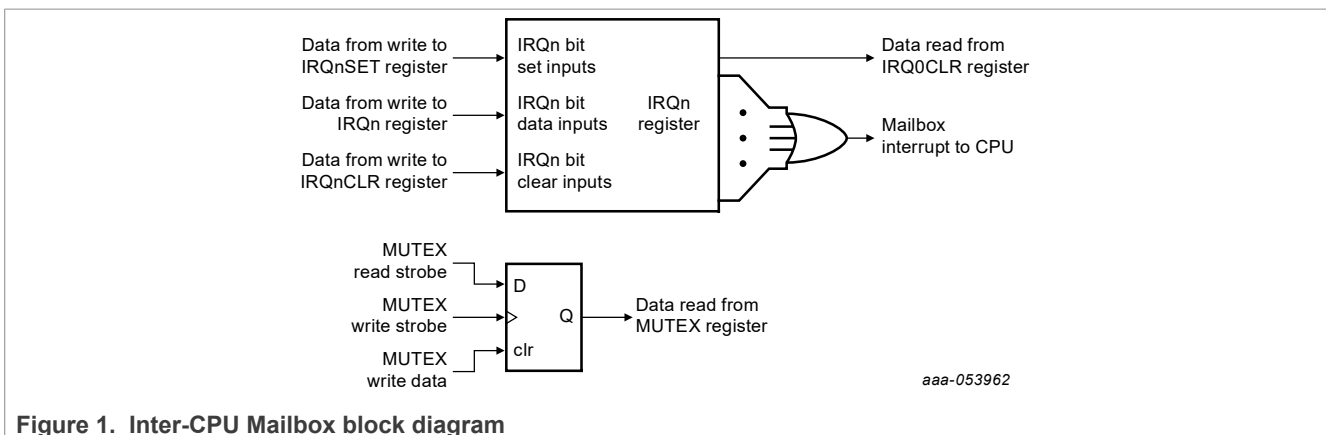


Figure 1. Inter-CPU Mailbox block diagram

For the following two examples, one can consider that the CPU1 is already booted. By default, the CPU1 image is loaded and executed from the FLASH memory, however there is the possibility to boot the CPU1 image from RAM by defining the preprocessor symbol CORE1_IMAGE_COPY_TO_RAM to 1 in the CPU0 project settings. The two examples below use the MCUXpresso SDK 2.13.1.

Also, the initial configuration of the MAILBOX can be accomplished as follow, whatever the mechanism used:

- Enable the clock to the MAILBOX module by writing a 1 to bit 31 in the SYSCON AHBCLKCTRLSET register.
- Clear the peripheral reset in the SYSCON register by writing bit 31 (MAILBOX) to the PRESETCTRLCLR0 register.
- Enable the interrupt in the NVIC to handle it when the core receives an interrupt in normal mode.

3.1 Interrupt mechanism

This section provides information about how to configure and use the MAILBOX for inter-core communication using the interrupt mechanism.

3.1.1 Example

Once the MAILBOX is properly configured using the above instruction, CPU0 and CPU1 can communicate by sending messages that generate automatically an interrupt to the other core. Messages are contained in the 32-bits register IRQn and determined by the user based on requirements. Therefore, each bit of this register can be configured differently. Writing 1 to multiple bits of this register generates more than one interrupt.

• CPU0 sends message to CPU1

1. CPU0 writes a message to the secondary mailbox register by writing the message into the IRQ1SET register of the receiver core (CPU1).

Note: Writing into the IRQSET register writes 1 to the corresponding bit of IRQ.

1. An interrupt is generated in CPU1 that can be processed in the mailbox interrupt handler.
2. CPU1 reads the message in its IRQ1 register.
3. CPU1 clears the wanted bits in the IRQ1 register by writing the value into IRQ1CLR.

• CPU1 sends message to CPU0

1. CPU1 writes a message to the primary mailbox register by writing the message into the IRQ0SET register of the receiver core (CPU0).

Note: Writing into the IRQSET register writes 1 to the corresponding bit of IRQ.

1. An interrupt is generated in CPU0 that can be processed in the mailbox interrupt handler.
2. CPU0 reads the message in the IRQ0 register.
3. CPU0 clears the wanted bits in the IRQ0 register by writing the value into IRQ0CLR.

Figure 2 illustrates the example above. First, CPU0 sends a message (value) to CPU1 using 32 bits in IRQ1SET, then CPU1 sends back the message (value + 1) to CPU0.

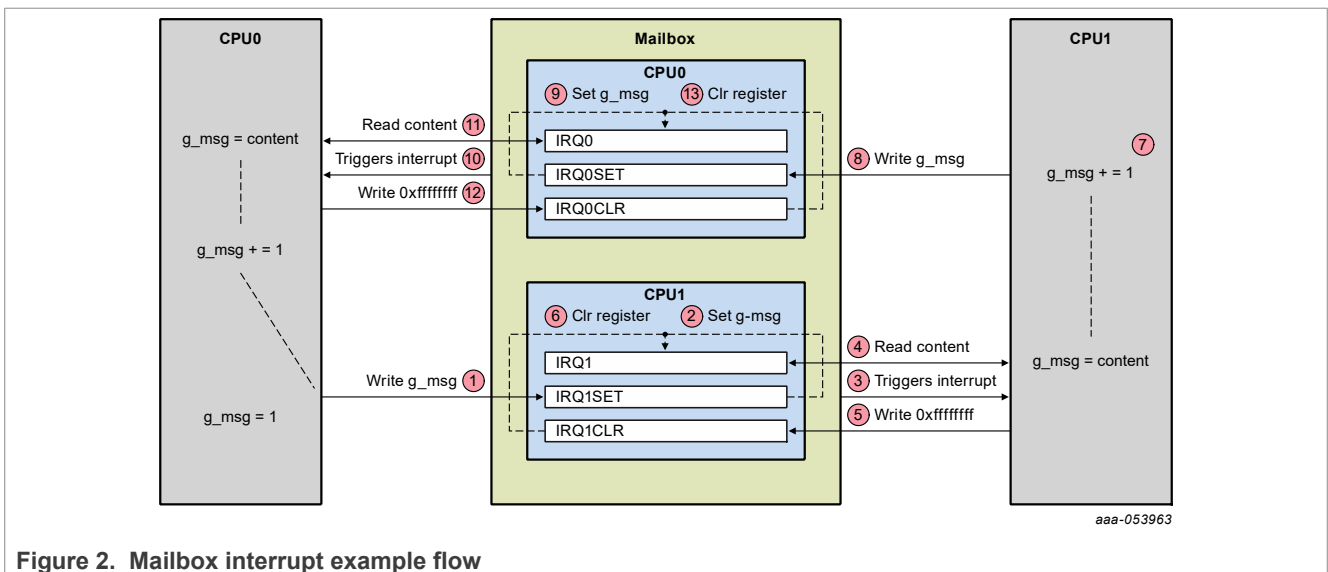


Figure 2. Mailbox interrupt example flow

3.1.2 Software implementation

MCUXpresso SDK for MCX Nx4x contains an example for using the interrupt mechanism of the Inter-CPU Mailbox at <SDK_LOCATION>\boards\<board_name>\driver_examples\mailbox\interrupt\. This SDK example shows the scenario explained above, where the primary core writes a value to the secondary core mailbox, it causes a mailbox interrupt on the secondary core side. The secondary core reads the value from the mailbox. It increments and writes it to the mailbox register for the primary core, which causes mailbox interrupt on the primary core side.

```
#define PRIMARY_CORE_MAILBOX_CPU_ID kMAILBOX_CM33_Core0
#define SECONDARY_CORE_MAILBOX_CPU_ID kMAILBOX_CM33_Core1

#define START_EVENT 1234
```

Figure 3. Mailboxes macros defined for the application and start event

```

55241 #define MAILBOX_BASE (0x400B2000u)
55242 /** Peripheral MAILBOX base pointer */
55243 #define MAILBOX ((MAILBOX_Type *)MAILBOX_BASE)
    
```

Figure 4. Mailbox base address

```

114 int main(void)
115 {
116     /* Init board hardware.*/
117     /* attach main clock divide to FLEXCOMM0 (debug console) */
118     CLOCK_AttachClk(BOARD_DEBUG_UART_CLK_ATTACH);
119
120     BOARD_InitBootPins();
121     BOARD_InitBootClocks();
122     BOARD_InitDebugConsole();
123
124     PRINTF("Mailbox interrupt example\r\n");
125
126     /* Init Mailbox */
127     MAILBOX_Init(MAILBOX);
128
129     /* Enable mailbox interrupt */
130     NVIC_EnableIRQ(MAILBOX_IRQn);
131
132 #ifdef CORE1_IMAGE_COPY_TO_RAM
133     /* Calculate size of the image */
134     uint32_t core1_image_size;
135     core1_image_size = get_core1_image_size();
136     PRINTF("Copy CORE1 image to address: 0x%x, size: %d\r\n", CORE1_BOOT_ADDRESS, core1_image_size);
137
138     /* Copy application from FLASH to RAM */
139     memcpy((void *)CORE1_BOOT_ADDRESS, (void *)CORE1_IMAGE_START, core1_image_size);
140 #endif
141
142     /* Start the secondary core */
143     start_secondary_core(CORE1_BOOT_ADDRESS);
144
145     /* Wait for start and initialization of secondary core */
146     while (!g_secondary_core_started)
147     ;
148
149     PRINTF("Write to the secondary core mailbox register: %d\r\n", g_msg);
150     /* Write g_msg to the secondary core mailbox register - it causes interrupt on the secondary core */
151     MAILBOX_SetValue(MAILBOX, SECONDARY_CORE_MAILBOX_CPU_ID, g_msg);
152
153     while (1)
154     {
155         __WFI();
156     }
157 }
158
    
```

Figure 5. Main function of CPU0

```

88 /* When the secondary core writes to the primary core mailbox register it causes call of this irq handler,
89    in which the received value is read, incremented and sent again to the secondary core */
90 void MAILBOX_IRQHandler()
91 {
92     if (!g_secondary_core_started)
93     {
94         if (START_EVENT == MAILBOX_GetValue(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID))
95         {
96             g_secondary_core_started = true;
97         }
98         MAILBOX_ClearValueBits(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID, 0xffffffff);
99     }
100     else
101     {
102         g_msg = MAILBOX_GetValue(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID);
103         MAILBOX_ClearValueBits(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID, 0xffffffff);
104         PRINTF("Read value from the primary core mailbox register: %d\r\n", g_msg);
105         g_msg++;
106         PRINTF("Write to the secondary core mailbox register: %d\r\n", g_msg);
107         MAILBOX_SetValue(MAILBOX, SECONDARY_CORE_MAILBOX_CPU_ID, g_msg);
108     }
109 }
    
```

Figure 6. Mailbox interrupt handler CPU0

```

45= int main(void)
46 {
47     /* Init board hardware.*/
48     /* set BOD VBAT level to 1.65V */
49     BOARD_InitBootPins();
50
51     /* Initialize Mailbox */
52     MAILBOX_Init(MAILBOX);
53
54     /* Enable mailbox interrupt */
55     NVIC_EnableIRQ(MAILBOX_IRQn);
56
57     /* Let the other side know the application is up and running */
58     MAILBOX_SetValue(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID, (uint32_t)START_EVENT);
59
60     while (1)
61     {
62         __WFI();
63     }
64 }

```

Figure 7. Main function of CPU1

```

34= void MAILBOX_IRQHandler()
35 {
36     g_msg = MAILBOX_GetValue(MAILBOX, SECONDARY_CORE_MAILBOX_CPU_ID);
37     MAILBOX_ClearValueBits(MAILBOX, SECONDARY_CORE_MAILBOX_CPU_ID, 0xffffffff);
38     g_msg++;
39     MAILBOX_SetValueBits(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID, g_msg);
40 }

```

Figure 8. Mailbox interrupt handler CPU1

[Figure 9](#) shows the result of running the project displayed in the terminal.

```

Mailbox interrupt example
Write to the secondary core mailbox register: 1
Read value from the primary core mailbox register: 2
Write to the secondary core mailbox register: 3
Read value from the primary core mailbox register: 4
Write to the secondary core mailbox register: 5
Read value from the primary core mailbox register: 6
Write to the secondary core mailbox register: 7
Read value from the primary core mailbox register: 8
Write to the secondary core mailbox register: 9
Read value from the primary core mailbox register: 10
Write to the secondary core mailbox register: 11
Read value from the primary core mailbox register: 12
Write to the secondary core mailbox register: 13
Read value from the primary core mailbox register: 14
Write to the secondary core mailbox register: 15

```

Figure 9. Inter-CPU mailbox interrupt example output on a serial terminal

3.2 Mutex mechanism

This section provides information about how to configure and use the MAILBOX for inter-core communication using the mutex mechanism.

3.2.1 Example

Once the MAILBOX is properly configured following the above instruction ([Inter-CPU Mailbox](#)), CPU0 and CPU1 can also communicate by using a mutex to share and modify a variable consecutively. CPU0 starts by sending the address of the shared variable to CPU1 by using the mechanism seen in the previous chapter.

Once that memory location address has been shared, both cores can try to get the mutex in a while loop. When the mutex is available, a core takes the mutex and updates the shared variable and then sets the mutex to release it and allow access from the other core to the shared variable.

3.2.2 Software implementation

MCUXpresso SDK for MCX Nx4x contains an example for using the interrupt mechanism of the Inter-CPU Mailbox at <SDK_LOCATION>\boards\<board_name>\driver_examples\mailbox\mutex\. This example shows the scenario explained above, where the primary core writes the shared variable address to the secondary core mailbox that causes a mailbox interrupt on the secondary core side to process the shared variable. Then in a while loop, both cores try to get the mutex. When a core gets the mutex, it updates the shared variable by incrementing its value by 1 and releases the mutex.

Figure x shows the flow of the mailbox mutex example. For details on the interrupt mechanism of this example, refer to the previous example.

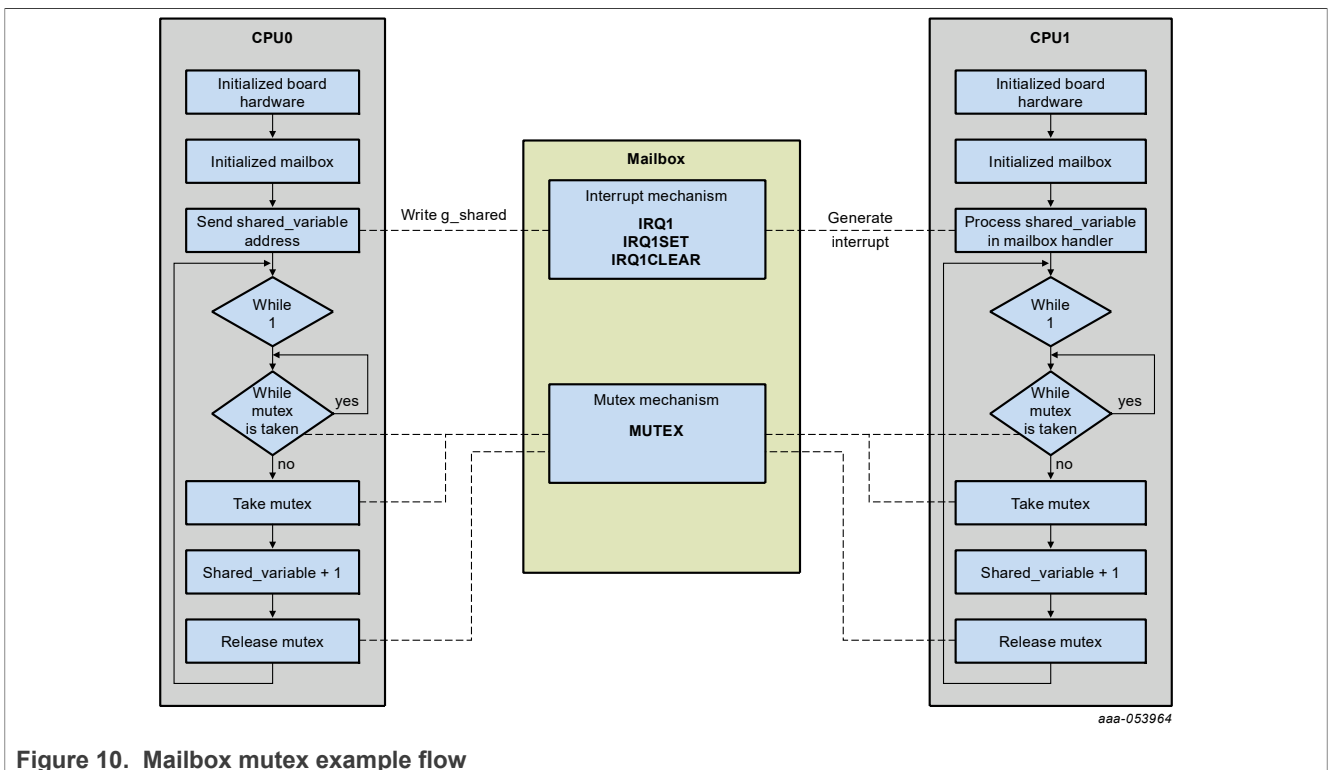


Figure 10. Mailbox mutex example flow

```

101=int main(void)
102 {
103     /* Init board hardware */
104     /* attach main clock divide to FLEXCOMM0 (debug console) */
105     CLOCK_AttachClk(BOARD_DEBUG_UART_CLK_ATTACH);
106
107     BOARD_InitBootPins();
108     BOARD_InitBootClocks();
109     BOARD_InitDebugConsole();
110
111     PRINTF("Mailbox mutex example\r\n");
112
113     /* Init Mailbox */
114     MAILBOX_Init(MAILBOX);
115
116     /* Enable mailbox interrupt */
117     NVIC_EnableIRQ(MAILBOX_IRQn);
118
119 #ifdef CORE1_IMAGE_COPY_TO_RAM
120     /* Calculate size of the image */
121     uint32_t core1_image_size;
122     core1_image_size = get_core1_image_size();
123     PRINTF("Copy CORE1 image to address: 0x%x, size: %d\r\n", CORE1_BOOT_ADDRESS, core1_image_size);
124
125     /* Copy application from FLASH to RAM */
126     memcpy((void *)CORE1_BOOT_ADDRESS, (void *)CORE1_IMAGE_START, core1_image_size);
127 #endif
128
129     /* Start the secondary core */
130     start_secondary_core(CORE1_BOOT_ADDRESS);
131
132     /* Wait for start and initialization of secondary core */
133     while (!g_secondary_core_started)
134         ;
135
136     /* Send address of shared variable to the secondary core */
137     MAILBOX_SetValue(MAILBOX, SECONDARY_CORE_MAILBOX_CPU_ID, (uint32_t)&g_shared);
138
139     while (1)
140     {
141         /* Get Mailbox mutex */
142         while (MAILBOX_GetMutex(MAILBOX) == 0)
143             ;
144
145         /* The core0 has mutex, can change shared variable g_shared */
146         g_shared++;
147
148         PRINTF("Core0 has mailbox mutex, update shared variable to: %d\r\n", g_shared);
149
150         /* Set mutex to allow access other core to shared variable */
151         MAILBOX_SetMutex(MAILBOX);
152
153         /* Add several nop instructions to allow the opposite core to get the mutex */
154         __asm("nop");
155         __asm("nop");
156         __asm("nop");
157         __asm("nop");
158         __asm("nop");
159         __asm("nop");
160         __asm("nop");
161         __asm("nop");
162         __asm("nop");
163         __asm("nop");
164     }
165 }
166

```

Figure 11. Main function of CPU0

```

89=void MAILBOX_IRQHandler()
90 {
91     if (START_EVENT == MAILBOX_GetValue(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID))
92     {
93         g_secondary_core_started = true;
94     }
95     MAILBOX_ClearValueBits(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID, 0xffffffff);
96 }

```

Figure 12. Mailbox interrupt handler of CPU0


```

49= int main(void)
50 {
51     /* Init board hardware */
52     /* set BOD VBAT level to 1.65V */
53     BOARD_InitBootPins();
54     BOARD_InitDebugConsole();
55
56     /* Initialize Mailbox */
57     MAILBOX_Init(MAILBOX);
58
59     /* Enable mailbox interrupt */
60     NVIC_EnableIRQ(MAILBOX_IRQn);
61
62     /* Let the other side know the application is up and running */
63     MAILBOX_SetValue(MAILBOX, PRIMARY_CORE_MAILBOX_CPU_ID, (uint32_t)START_EVENT);
64
65     while (1)
66     {
67         /* Get Mailbox mutex */
68         while (MAILBOX_GetMutex(MAILBOX) == 0)
69             ;
70
71         /* The core1 has mutex, can change shared variable g_shared */
72         if (g_shared != NULL)
73         {
74             (*g_shared)++;
75             PRINTF("Core1 has mailbox mutex, update shared variable to: %d\r\n", *g_shared);
76         }
77
78         /* Set mutex to allow access other core to shared variable */
79         MAILBOX_SetMutex(MAILBOX);
80
81         /* Add several nop instructions to allow the opposite core to get the mutex */
82         __asm("nop");
83         __asm("nop");
84         __asm("nop");
85         __asm("nop");
86         __asm("nop");
87         __asm("nop");
88         __asm("nop");
89         __asm("nop");
90         __asm("nop");
91         __asm("nop");
92     }
93 }
94

```

Figure 13. Main function of CPU1

```

31= /* Pointer to shared variable by both cores, before changing of this variable the
32    cores must first take Mailbox mutex, after changing the shared variable must
33    return mutex */
34 volatile uint32_t *g_shared = NULL;
35
36= /******
37    * Code
38    *****/
39= void MAILBOX_IRQHandler()
40 {
41     g_shared = (uint32_t *)MAILBOX_GetValue(MAILBOX, SECONDARY_CORE_MAILBOX_CPU_ID);
42     MAILBOX_ClearValueBits(MAILBOX, SECONDARY_CORE_MAILBOX_CPU_ID, 0xffffffff);
43     __DSB();
44 }

```

Figure 14. Mailbox interrupt handler of CPU1

```
Core1 has mailbox mutex, update shared variable to: 1
Core0 has mailbox mutex, update shared variable to: 2
Core1 has mailbox mutex, update shared variable to: 3
Core0 has mailbox mutex, update shared variable to: 4
Core1 has mailbox mutex, update shared variable to: 5
Core0 has mailbox mutex, update shared variable to: 6
Core1 has mailbox mutex, update shared variable to: 7
Core0 has mailbox mutex, update shared variable to: 8
Core1 has mailbox mutex, update shared variable to: 9
Core0 has mailbox mutex, update shared variable to: 10
Core1 has mailbox mutex, update shared variable to: 11
Core0 has mailbox mutex, update shared variable to: 12
Core1 has mailbox mutex, update shared variable to: 13
Core0 has mailbox mutex, update shared variable to: 14
Core1 has mailbox mutex, update shared variable to: 15
Core0 has mailbox mutex, update shared variable to: 16
Core1 has mailbox mutex, update shared variable to: 17
Core0 has mailbox mutex, update shared variable to: 18
Core1 has mailbox mutex, update shared variable to: 19
Core0 has mailbox mutex, update shared variable to: 20
```

Figure 15. Inter-CPU mailbox mutex example output on a serial terminal

4 Conclusion

This application note provides a quick look at the dual-core mechanisms in the MCX Nx4x device. It briefly introduces some applications related to dual-core communication based on the MAILBOX driver examples in the MCX Nx4x SDK and demonstrates how to use the interrupt and mutex mechanisms depending on the use case.

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6 Revision history

Table 1. Revision history

Document ID	Release date	Description
AN14150 v.1.0	20 January 2024	Initial version

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