

Semiconductor Products Sector
Application Note

AN2153

A Serial Bootloader for Reprogramming the MC9S12DP256 FLASH Memory

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Introduction

The MC9S12DP256 is a member of the M68HC12 Family of 16-bit microcontrollers (MCU) containing 262,144 bytes of bulk or sector erasable, word programmable FLASH memory arranged as four 65,536 byte blocks. Including FLASH memory, rather than EPROM or ROM, on a microcontroller has significant advantages.

For the manufacturer, placing system firmware in FLASH memory provides several benefits. First, firmware development can be extended late into the product development cycle by eliminating masked ROM lead times. Second, when a manufacturer has several products based on the same microcontroller, it can help eliminate inventory problems and lead times associated with ROM-based microcontrollers. Finally, if a severe bug is found in the product's firmware during the manufacturing process, the in-circuit reprogrammability of FLASH memory prevents the manufacturer from having to scrap any work-in-process.

The ability of FLASH memory to be electrically erased and reprogrammed also provides benefits for the manufacturer's end customers. The customer's products can be updated or enhanced with new features and capabilities without having to replace any components or return a product to the factory.

Unlike the M68HC11 Family, the MC9S12DP256 does not have a bootstrap ROM containing firmware to allow initial programming of the FLASH directly through one of the on-chip serial communications interface (SCI) ports. Initial on-chip FLASH programming requires either special test and handling equipment to program the device before it is placed in the target system or a background debug module (BDM) programming tool available from Freescale or a third party vendor.

The MC9S12DP256's four on-chip FLASH arrays contain two variable size, erase protectable areas as shown in [Figure 1](#). While the majority of the bootloader could be contained in any of the protected areas, the protected high area in the \$C000–\$FFFF memory range must at least contain reset and interrupt vectors that point to a jump table. In most cases, unless a complex or sophisticated communication protocol is required that will not fit into 16 K, it is easiest to place the entire bootloader into the protected high area of block zero.

Erasing and programming the on-chip FLASH memory of the MC9S12DP256 presents some unique challenges. Even though FLASH block zero has two separate erase protected areas, code cannot be run out of either protected area while the remainder of the block is erased or programmed. While it is possible to run code from one FLASH block while erasing or reprogramming another, adopting such a strategy would complicate the overall implementation of the bootloader. Consequently, during the erase and reprogram process, the code must reside in other on-chip memory or in external memory. In addition, because the reset and interrupt vectors reside in the erase protected area, they cannot be changed. This necessitates a secondary reset/interrupt vector table be placed outside the protected FLASH memory area.

The remainder of this application note explores the requirements of a serial bootloader and the implementation of the programming algorithm for the MC9S12DP256's FLASH.

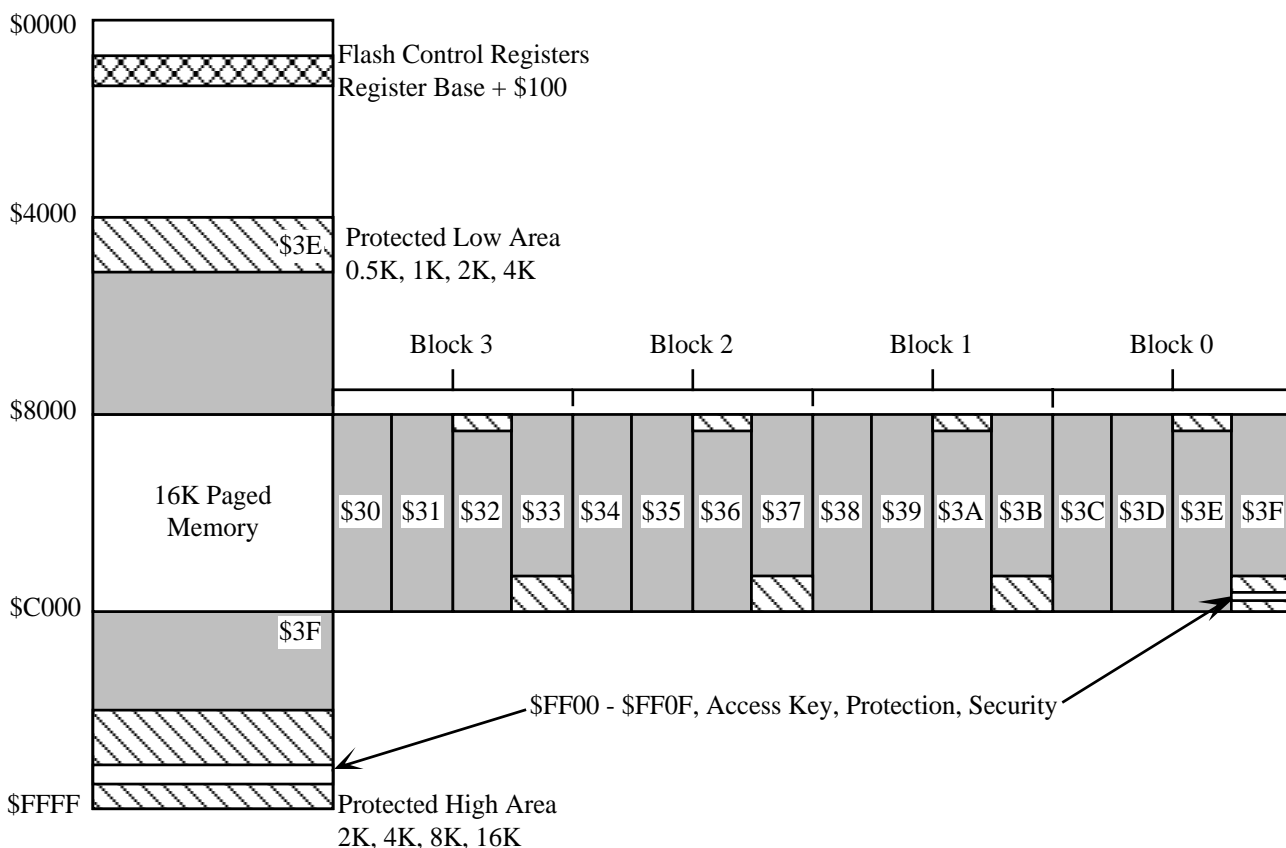


Figure 1. MC9S12DP256 Memory Map

Overview of the MC9S12DP256's FLASH

The MC9S12DP256's 256 K of on-chip FLASH memory is composed of four 65,536 byte blocks. Each block is arranged as 32,768 16-bit words and may be read as bytes, words, or misaligned words. Access time is one bus cycle for bytes and aligned words reads and two bus cycles for misaligned word reads. Write operations for program and erase operations can be performed only as an aligned word. Each 64-K block is organized in 1024 rows of 32 words. An erase sector contains 8 rows or 512 bytes. Erase operations may be performed on a sector as small as 512 bytes or on the entire 65,536-byte block. An erased word reads \$FFFF and a programmed word reads \$0000.

The programming voltage required to program and erase the FLASH is generated internally by on-chip charge pumps. Program and erase operations are performed by a command driven interface from the microcontroller using an internal state machine. The completion of a program or erase operation is signaled by the setting of the CCIF flag and may optionally generate an interrupt. All FLASH blocks can be programmed or erased at the same time; however, it is not possible to read from a FLASH block while it is being erased or programmed.

Each 64-K block contains hardware interlocks which protect data from accidental corruption. As shown in **Figure 1**, the upper 32 K of block zero can be accessed through the 16-Kbyte PPAGE window or at two fixed address 16-K address ranges. One protected area is located in the upper address area of the fixed page address range from \$C000–\$FFFF and is normally used for bootloader code. Another area is located in the lower portion of the fixed page address range from \$4000–\$7FFF. Additional protected memory areas are present in the three remaining 64-K FLASH blocks; however, they are only accessible through the 16-K PPAGE window.

FLASH Control Registers

The control and status registers for all four FLASH blocks occupy 16 bytes in the input/output (I/O) register area. To accommodate the four FLASH blocks while occupying a minimum of register address space, the FLASH control register address range is divided into two sections. The first four registers, as shown in **Figure 2**, apply to all four memory blocks. The remaining 12 bytes of the register space have duplicate sets of registers, one for each FLASH bank. The active register bank is selected by the BKSEL bits in the unbanked FLASH configuration register (FCNFG). Note that only three of the banked registers contain usable status and control bits; the remaining nine registers are reserved for factory testing or are unused.

	Bit 7	6	5	4	3	2	1	Bit 0	
FCLKDIV	FDIVLD	PRDIV8	FDIV5	FDIV4	FDIV3	FDIV2	FDIV1	FDIV0	\$x100
FSEC	KEYEN	NV6	NV5	NV4	NV3	NV2	SEC01	SEC00	\$x101
Reserved	0	0	0	0	0	0	0	0	\$x102
FCNFG	CBEIE	CCIE	KEYACC	0	0	0	BKSEL1	BKSEL1	\$X103
Unbanked									
Banked									
FPROT	FPOPEN	F	FPHDIS	FPHS1	FPHS0	FPLDIS	FPLS1	FPLS0	\$X104
FSTAT	CBEIF	CCIF	PVIOL	ACCERR	0	BLANK	0	0	\$X105
FCMD	0	ERASE	PROG	0	0	ERVER	0	MASS	\$X106
Reserved	0	0	0	0	0	0	0	0	\$X107– \$x10F

Figure 2. FLASH Status and Control Registers

FLASH Protection

The protected areas of each FLASH block are controlled by four bytes of FLASH memory residing in the fixed page memory area from \$FF0A–\$FF0D. During the microcontroller reset sequence, each of the four banked FLASH protection registers (FPROT) is loaded from values programmed into these memory locations. As shown in [Figure 3](#), location \$FF0A controls protection for block three, \$FF0B controls protection for block two, \$FF0C controls protection for block one, and \$FF0D controls protection for block zero.

The values loaded into each FPROT register determine whether the entire block or just subsections are protected from being accidentally erased or programmed. As mentioned previously, each 64-K block can have two protected areas. One of these areas, known as the lower protected block, grows from the middle of the 64-K block upward. The other, known as the upper protected block, grows from the top of the 64-K block downward. In general, the upper protected area of FLASH block zero is used to hold bootloader code since it contains the reset and interrupt vectors. The lower protected area of block zero and the protected areas of the other FLASH blocks can be used for critical parameters that would not change when program firmware was updated.

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The FPOPEN bit in each FPROT register determines whether the entire FLASH block or subsections of it can be programmed or erased. When the FPOPEN bit is erased (1), the remainder of the bits in the register determine the state of protection and the size of each protected block. In its programmed state (0), the entire FLASH block is protected and the state of the remaining bits within the FPROT register is irrelevant.

Address	Description
\$FF00–\$FF07	Security back door comparison key
\$FF08–\$FF09	Reserved
\$FF0A	Protection byte for FLASH block 3
\$FF0B	Protection byte for FLASH block 2
\$FF0C	Protection byte for FLASH block 1
\$FF0D	Protection byte for FLASH block 0
\$FF0E	Reserved
\$FF0F	Security byte

Figure 3. FLASH Protection and Security Memory Locations

The FPHDIS and FPLDIS bits determine the protection state of the upper and lower areas within each 64-K block respectively. The erased state of these bits allows erasure and programming of the two protected areas and renders the state of the FPHS[1:0] and FPLS[1:0] bits immaterial. When either of these bits is programmed, the FPHS[1:0] and FPLS[1:0] bits determine the size of the upper and lower protected areas. The tables in **Figure 4** summarize the combinations of the FPHS[1:0] and FPLS[1:0] bits and the size of the protected area selected by each.

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FPHS[1:0]	Protected Size	FPLS[1:0]	Protected Size
0:0	2 K	0:0	512 bytes
0:1	4 K	0:1	1 K
1:0	8 K	1:0	2 K
1:1	16 K	1:1	4 K

Figure 4. FLASH Protection Select Bits

The FLASH protection registers are loaded during the reset sequence from address \$FF0D for FLASH block 0, \$FF0C for FLASH block 1, \$FF0B for FLASH block 2 and \$FF0A for FLASH block 3. This is indicated by the “F” in the reset row of the register diagram in the MC9S12DP256 data book. This register determines whether a whole block or subsections of a block are protected against accidental program or erase. Each FLASH block can have two protected areas, one starting from relative address \$8000 (called lower) toward higher addresses and the other growing downward from \$FFFF (called higher). While the later is mainly targeted to hold the bootloader code since it covers the vector space (FLASH 0), the other area may be used to keep critical parameters. Trying to alter any of the protected areas will result in a protect violation error, and bit PVIOL will be set in the FLASH status register FSTAT.

NOTE: *A mass or bulk erase of the full 64-Kbyte block is only possible when the FPLDIS and FPHDIS bits are in the erased state.*

FLASH Security

The security of a microcontroller’s program and data memories has long been a concern of companies for one main reason. Because of the considerable time and money that is invested in the development of proprietary algorithms and firmware, it is extremely desirable to keep the firmware and associated data from prying eyes. This was an especially difficult problem for earlier M68HC12 Family members as the background debug module (BDM) interface provided easy, uninhibited access to the FLASH and EEPROM contents using a 2-wire connection. Later revisions of the original D Family parts provided a method that

allowed a customer's firmware to disable the BDM interface (BDM lockout) once the part had been placed in the circuit and programmed. While this prevents the FLASH and EEPROM from being easily accessed in-circuit, it does not prevent a D Family part from being removed from the circuit and placed in expanded mode so the FLASH and EEPROM can be read.

The security features of the MC9S12DP256 have been greatly enhanced. While no security feature can be 100 percent guaranteed to prevent access to an MCU's internal resources, the MC9S12DP256's security mechanism makes it extremely difficult to access the FLASH or EEPROM contents. Once the security mechanism has been enabled, access to the FLASH and EEPROM either through the BDM or the expanded bus is inhibited. Gaining access to either of these resources may be accomplished only by erasing the contents of the FLASH and EEPROM or through a built-in back door mechanism. While having a back door mechanism may seem to be a weakness of the security mechanism, the target application must specifically support this feature for it to operate.

Erasing the FLASH or EEPROM can be accomplished using one of two methods. The first method requires resetting the target MCU in special single-chip mode and using the BDM interface. When a secured device is reset in special single-chip mode, a special BDM security ROM becomes active. The program in this small ROM performs a blank check of the FLASH and EEPROM memories. If both memory spaces are erased, the BDM firmware temporarily disables device security, allowing full BDM functionality. However, if the FLASH or EEPROM are not blank, security remains active and only the BDM hardware commands remain functional. In this mode, the BDM commands are restricted to reading and writing the I/O register space. Because all other BDM commands and on-chip resources are disabled, the contents of the FLASH and EEPROM remain protected. This functionality is adequate to manipulate the FLASH and EEPROM control registers to erase their contents.

NOTE: *Use of the BDM interface to erase the FLASH and EEPROM memories is not present in the initial mask set (OK36N) of the MC9S12DP256. Great care must be exercised to ensure that the microcontroller is not programmed in a secure state unless the back door mechanism is supported by the target firmware.*

The second method requires the microcontroller to be connected to external memory devices and reset in expanded mode where a program can be executed from the external memory to erase the FLASH and EEPROM. This method may be preferred before parts are placed in a target system.

As shown in **Figure 5**, the security mechanism is controlled by the two least significant bits in the security byte. Because the only unsecured combination is when SEC1 has a value of 1 and SEC0 has a value of 0, the microcontroller will remain secured even after the FLASH and EEPROM are erased, since the erased state of the security byte is \$FF. As previously explained, even though the device is secured after being erased, the part may be reset in special single-chip mode, allowing manipulation of the microcontroller via the BDM interface. However, after erasing the FLASH and EEPROM, the microcontroller can be placed in the unsecured state by programming the security byte with a value of \$FE. Note that because the FLASH must be programmed one aligned word at a time and because the security byte resides at an odd address (\$FF0F), the word at \$FF0E must be programmed with a value of \$FFFE.

SEC[1:0]	Security State
0:0	Secured
0:1	Secured
1:0	Unsecured
1:1	Secured

Figure 5. Security Bits

**Utilizing the
 FLASH Security
 Back Door**

In normal single-chip or normal expanded operating modes, the security mechanism may be temporarily disabled only through the use of the back door key access feature. Because the back door mechanism requires support by the target firmware, it is impossible for the back door mechanism to be used to defeat device security unless the capability is designed into the target application. To disable security, the firmware must have access to the 64-bit value stored in the security back door comparison key located in FLASH memory from \$FF00–\$FF07. If

operating in single-chip mode, the key would typically be provided to the firmware through one of the on-chip serial ports. In addition, back door security bypass must be enabled by leaving the most significant bit of the Security byte at \$FF0F erased. To disable the back door security bypass feature, this bit should be programmed to zero.

Once the application receives the 64-bit key, it must set the KEYACC bit in the FCNFG register. After setting the KEYACC bit, the firmware must write the received 64-bit key to the security back door comparison key memory locations (\$FF00–\$FF07) as four 16-bit words, in sequential order. Finally, the KEYACC bit must be cleared. If all four 16-bit words written to the comparison key memory area matched the corresponding values stored in FLASH, the MCU will be unsecured by forcing the SEC[1:0] bits in the FSEC register to the unsecured state. Note that this operation only temporarily disables the device security. The next time the MCU is reset, the SEC[1:0] bits will be loaded from the security byte at \$FF0F

FLASH Program and Erase Overview

All FLASH program and erase timings are handled by a hardware state machine, freeing the CPU to perform other tasks during these operations. The timebase for the state machine is derived from the oscillator clock via a programmable down counter. Program and erase operations are accomplished by writing values to the FCMD register. Four commands are recognized in the current implementation and are summarized in [Figure 6](#).

Command	Operation	Description
\$20	Memory program	Program 1 aligned word, 2 bytes
\$40	Sector erase	Erase a 512-byte sector
\$41	Mass erase	Erase a 64-Kbyte block
\$05	Erase verify	Verify erasure of a 64-Kbyte block
Other	Illegal	Generate an access error

Figure 6. FLASH Program and Erase Commands

The command register and the associated address and data registers are implemented as a 2-stage first in, first out (FIFO) command buffer. This configuration allows a new command to be issued while the hardware state machine completes the previously issued command. The main reason for this design is to decrease programming time. Without the 2-stage FIFO command buffer, the programming voltage would have to be removed from the FLASH array at the end of each program command to avoid exceeding the high voltage active time, t_{HV} , specification. Applying and removing the programming voltage after each program command would double the time required to program an aligned word. If program commands are continuously available to the state machine, it will keep high voltage applied to the array if the program command operates on the same 64-byte row. If the command in the second stage of the FIFO buffer has changed, the address is not within the same 64-byte row or the command buffer is empty, the high voltage will be removed and reapplied with a new command if required.

To aid the development of a multitasking environment where the CPU can perform other tasks while performing program and erase operations, the FLASH module control registers provide the ability to generate interrupts when a command completes or the command buffer is empty. When the command buffers empty interrupt enable (CBEIE) bit is set, an interrupt is generated whenever the command buffers empty interrupt flag (CBEIF) is set. When the command complete interrupt enable (CCIE) bit is set, an interrupt is generated when the command complete interrupt flag (CCIF) is set. Note that the CCIF flag is set at the completion of each command while the CBEIF is set when both stages of the FIFO are empty.

NOTE: *Because the interrupt vectors are located in FLASH block zero, memory locations in block zero cannot be erased or programmed when utilizing FLASH interrupts in a target application.*

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FLASH Erasure

As previously discussed, each 64-K block is organized in 1024 rows of 32 words. An erase sector contains 8 rows or 512 bytes. Erase operations may be performed on a sector as small as 512 bytes or on the entire 65,536 byte block. An erased word reads \$FFFF and a programmed word reads \$0000. Program and erase operations are very similar, differing only in the command written to the FCMD register and the data written to the FLASH memory array. The FLASH state machine erase and verify command operation is depicted in the flowchart of **Figure 7**.

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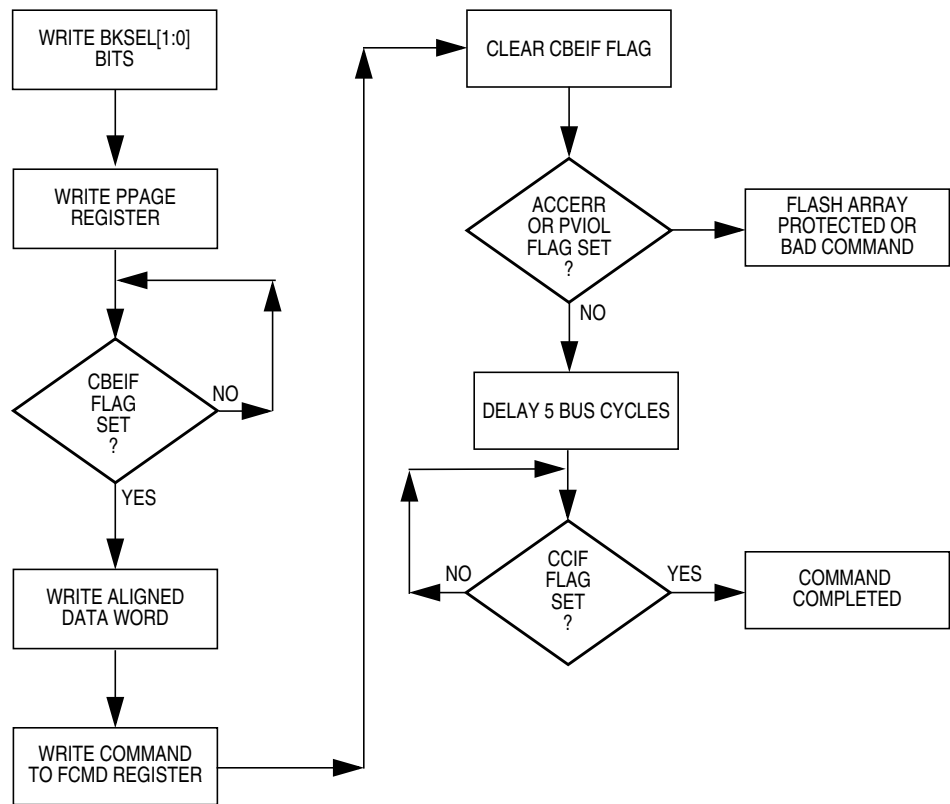


Figure 7. Erase and Verify Flowchart

Before beginning either an erase or program operation, it is necessary to write a value to the FCLKDIV register. The value written to the FCLKDIV register programs a down counter used to divide the oscillator clock, producing a 150-kHz to 200-kHz clock source used to drive the FLASH memory’s state machine. The most significant bit of the FCLKDIV register, when set, indicates that the register has been

initialized. If FDIVLD is clear, it indicates that the register has not been written to since the part was last reset. Attempting to erase or program the FLASH without initializing the FCLKDIV register will result in an access error and the command will not be executed.

A combination of the PRDIV8 and FDIV[5:0] bits is used to divide the oscillator clock to the 150-kHz to 200-kHz range required by the FLASH's state machine. The PRDIV8 bit is used to control a 3-bit prescaler. When set, the oscillator clock will be divided by eight before being fed to the 6-bit programmable down counter. Note that if the oscillator clock is greater than 12.8 MHz, the PRDIV8 bit must be set to obtain a proper state machine clock source using the FDIV[5:0] bits. The formulas for determining the proper value for the FDIV[5:0] bits are shown in **Figure 8**.

```

if (OSCCLK > 12.8 MHz)
    PRDIV8 = 1
else
    PRDIV8 = 0

if (PRDIV8 == 1)
    CLK = OSCCLK / 8
else
    CLK = OSCCLK

FCLKDIV[5:0] = INT((CLK / 1000) / 200)

FCLK = CLK / (FCLKDIV[5:0] + 1)

```

Figure 8. FCLKDIV Formulas

In the formulas, OSCCLK represents the reference frequency present at the EXTAL pin, NOT the bus frequency or the PLL output. The INT function always rounds toward zero and FCLK represents the frequency of the clock signal that drives the FLASH's state machine.

NOTE: *Erasing or programming the FLASH with an oscillator clock less than 500 kHz should be avoided. Setting FCLKDIV such that the state machine clock is less than 150 kHz can destroy the FLASH due to high voltage over stress. Setting FCLKDIV such that the state machine clock is greater than 200 kHz can result in improperly programmed memory locations.*

After initializing the FCLKDIV register with the proper value, the PPAGE register and the BKSEL[1:0] bits must be initialized. The PPAGE register must be written with a value that places the correct 16-K memory block in the PPAGE window that contains the memory area to be erased. If a mass (bulk) erase operation is performed on one of the 64-K blocks, the PPAGE register may be written with any one of the four PPAGE values associated with a 64-K block. Note that when performing a mass or sector erase in the address range of one of the two fixed pages, \$4000–\$7FFF or \$C000–\$FFFF, the value of the PPAGE register is unimportant.

The BKSEL[1:0] bits, located in the FCNFG register, are used to select the banked status and control registers associated with the 64-K FLASH block in which the erase operation is to be performed. As shown in **Figure 1**, the value of the FLASH block number decreases with increasing PPAGE values. Closely examining **Figure 1** reveals that the correct value for the BKSEL[1:0] bits is the one's complement of the PPAGE[3:2] register bits. Even though the flowchart shows the block select bits being written before the PPAGE register, these registers may be written in reverse order. This makes the code implementation straight forward since the value of the block select bits may be easily derived from the value written to the PPAGE register.

After initializing the PPAGE register and the block select bits, the command buffer empty interrupt flag (CBEIF) bit should be checked to ensure that the address, data and command buffers are empty. If the CBEIF bit is set, the buffers are empty and a program or erase command sequence can be started. The next three steps in the flowchart must be strictly adhered to. Any intermediate writes to the FLASH control and status registers or reads of the FLASH block on which the operation is being performed will cause the access error (ACCERR) flag to be set and the operation will be immediately terminated. For a mass erase operation, the address of the aligned data word may be any valid address in the 64-K block. For a sector erase, only the upper seven address bits are significant, the lower eight bits are ignored. For all erase operations, the data written to the FLASH block is ignored.

After writing a program or erase command to the FCMD register, the CBEIF bit must be written with a value of 1 to clear the CBEIF bit and initiate the command. After clearing the CBEIF bit, the ACCERR and PVIOL bits should be checked to ensure that the command sequence was valid. If either of these bits is set, it indicates that an erroneous command sequence was issued and the command sequence will be immediately terminated. Note that if either or both of the ACCERR and PVIOL bits are set, they must be cleared by writing a 1 to each flag's associated bit position before another command sequence can be initiated. Five bus cycles after the CBEIF bit is cleared, the CCIF flag will be cleared by the state machine indicating that the command was successfully begun. If a previous command has not been issued, the CBEIF bit will become set, indicating that the address, data, and command buffers are available to begin a new command sequence.

Once the erase command has completed, erasure of the sector or block should be verified to ensure that all locations contain \$FF. When erasing a 512-byte sector, each byte or word must be checked for an erased condition using software. Fortunately, however, the state machine has a verify command built into the hardware to perform an erase verify on the contents of any of the 64-K blocks. The command sequence used to perform an erase verify is identical to that of performing an erase command except that the erase verify command (\$05) is written to the FCMD register and the block select bits and the PPAGE register need not be rewritten. If all locations in a 64-K block are erased, a successful erase verify will cause the BLANK bit in the FSTAT register to be set. Note that the BLANK bit must be cleared by writing a 1 to its associated bit position before the next erase verify command is issued.

FLASH Programming

As mentioned in the previous section, the erase and program operations follow a nearly identical flow. There are, however, some minor changes to the flow that can improve the efficiency of the programming process. To take advantage of the decreased programming time provided by the 2-stage FIFO command buffer, it must be kept full with programming commands. As the flowchart in [Figure 9](#) shows, rather than waiting for each programming command to complete, a new programming command is issued as soon as the CBIEF flag is set. This allows the programming voltage to remain applied to the array as long as the next

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aligned word address remains within the same 64-byte row. Therefore, to minimize programming times, blocks of data to be programmed into the FLASH array should begin on a 64-byte boundary and be a multiple of 64 bytes.

Verification of programmed data should be performed only after a block of data has been programmed and all programming commands have completed. Performing a read operation on the FLASH array while a programming command is executing will cause the ACCERR flag to be set and all current and pending commands are terminated.

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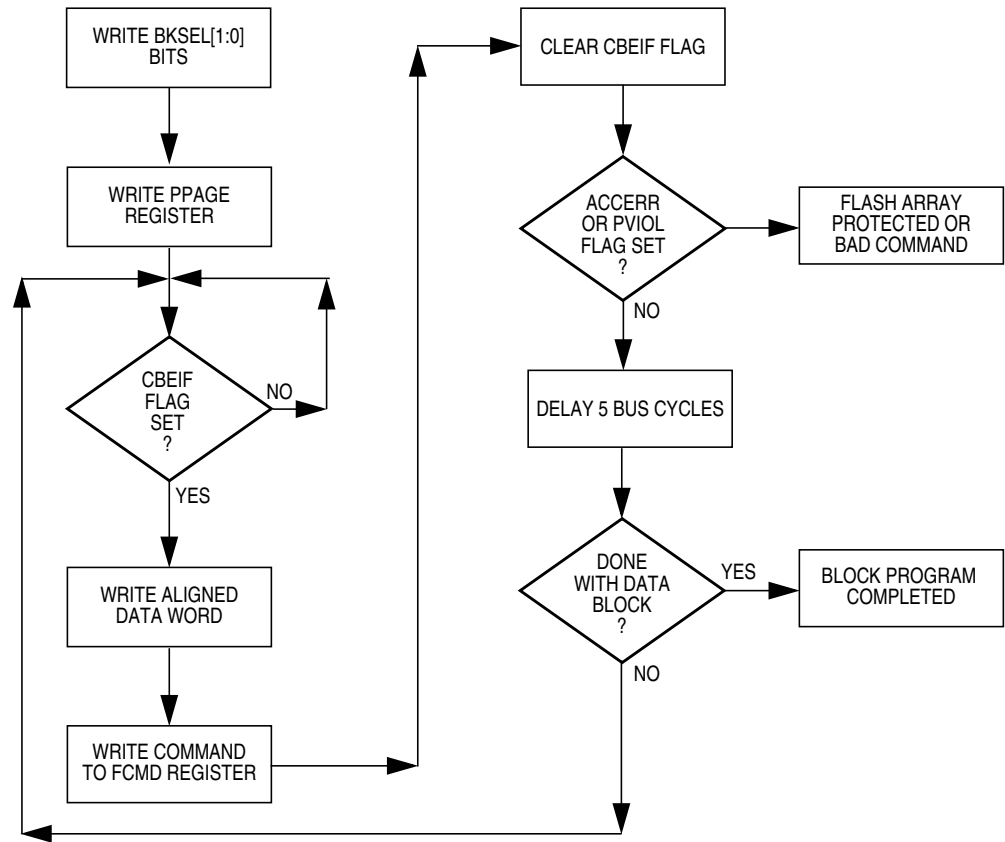


Figure 9. Programming Flowchart

General FLASH Serial Bootloader Requirements

A program such as the FLASH serial bootloader has two important requirements. First, it must have minimal impact on the final product's software performance. Second, it should add little or no cost to the hardware design. Because the MC9S12DP256 includes a variety of on-chip communications modules, five CAN modules, one J1850 module, two SCI ports, and three SPI modules, no additional external hardware should be required. Designs incorporating a CAN or J1850 network connection could easily incorporate the existing connection into the bootloader to download the new FLASH data. For applications not utilizing a network connection in the basic design, one of the two SCI ports can be used. In many systems, the SCI may be a part of the hardware design since it is often used as a diagnostic port. If an RS232 level translator is not included as part of the system design, a small adapter board can be constructed containing the level translator and RS232 connector. This board can then be used by service personnel to update the system firmware. Using such an adapter board prevents the cost of the level translator and connector from being added to each system. In addition to the SCI port, a single input pin is required to notify the serial bootloader startup code to execute the bootloader code or jump to the system application program.

As mentioned previously, because the MC9S12DP256's interrupt and reset vectors reside in the protected bootblock, they cannot be changed without erasing the bootblock itself. Even though it is possible to erase and reprogram the bootblock, it is inadvisable to do so. If anything goes wrong during the process of reprogramming the bootblock, it would be impossible to recover from the situation without the use of BDM programming hardware. For this reason, a bootloader should include support for a secondary interrupt and reset vector table located just below the protected bootblock area. Each entry in the secondary interrupt table should consist of a 2-byte address mirroring the primary interrupt and reset vector table. The secondary interrupt and reset vector table is utilized by having each vector point to a single JMP instruction that uses the CPU12's indexed-indirect program counter relative addressing mode. This form of the JMP instruction uses four bytes of

memory and requires just six CPU clock cycles to execute. For systems operating at the maximum bus speed of 25.0 MHz, six bus cycles adds only 240 ns to the interrupt latency. In most applications, this small amount of additional time will not affect the overall performance of the system.

Bootloader S-Record Format

The S-record object file format was designed to allow binary object code and/or data to be represented in printable ASCII hexadecimal format allowing easy transportation between computer systems and development tools. For M68HC12 Family members supporting less than 64 Kbytes of address space, S1 records, which contain a 16-bit address, are sufficient to specify the location in the device's memory space where code and/or data are to be loaded. The load address contained in the S1 record generally corresponds directly to the address of on-chip or off-chip memory device. For M68HC12 devices that support an address space greater than 64 Kbytes, S1 records are not sufficient.

Because the M68HC12 Family is a 16-bit microcontroller with a 16-bit program counter, it cannot directly address a total of more than 64 Kbytes of memory. To enable the M68HC12 Family to address more than 64 Kbytes of program memory, a paging mechanism was designed into the architecture. Program memory space expansion provides a window of 16-Kbyte pages that are located from \$8000–\$BFFF. An 8-bit paging register, called the PPAGE register, provides access to a maximum of 256, 16-Kbyte pages or 4 megabytes of program memory. While there may never be any devices that contain this much on-chip memory, the MC68HC812A4 is capable of addressing this much external memory. In addition, the MC9S12DP256 contains 256 Kbytes of on-chip FLASH residing in a 1MB address space.

While many high-level debuggers are capable of directly loading linked, absolute binary object files into a target system's memory, the bootloader does not have that ability. The bootloader is only capable of loading object files that are represented in the S-record format. Because S1 records only contain a 16-bit address, they are inadequate to specify a load address for a memory space greater than 64 Kbytes. S2 records, which contain a 24-bit load address, were originally defined for loading object files into the memory space of the M68000 Family. It would seem

that S2 records would provide the necessary load address information required for M68HC12 object files. However, as those who are familiar with the M68000 Family know, the M68000 has a linear (non-paged) address space. Thus, development tools, such as non-volatile memory device programmers, interpret the 24-bit address as a simple linear address when placing program data into memory devices.

Because the M68HC12 memory space expansion is based on 16-Kbyte pages, there is not a direct one-to-one mapping of the 24-bit linear address contained in the S2 record to the 16-Kbyte program memory expansion space. Instead of defining a new S-record type or utilizing an existing S-record type in a non-standard manner, the bootloader's program FLASH command views the MC9S12DP256's memory space as a simple linear array of memory that begins at an address of \$C0000. This is the same format in which S-records would need to be presented to a stand alone non-volatile memory device programmer.

The MC9S12DP256 implements six bits of the PPAGE register which gives it a 1MB program memory address space that is accessed through the PPAGE window at addresses \$8000–\$BFFF. The lower 768-K portion of the address space, accessed with PPAGE values \$00–\$2F, are reserved for external memory when the part is operated in expanded mode. The upper 256 K of the address space, accessed with PPAGE values \$30–\$3F, is occupied by the on-chip FLASH memory. The mapping between the linear address contained in the S-record and the 16-Kbyte page viewable through the PPAGE is shown in [Figure 10](#).

The generation of S-records that meet these requirements is the responsibility of the linker and/or S-record generation utility provided by the compiler/assembler vendor. Cosmic Software's linker and S-record generation utility is capable of producing properly formatted S-records that can be used by the bootloader. Other vendor's tools may or may not possess this capability. For those compilers and assemblers that produce "banked" S-records, an S-record conversion utility, SRecCvt.exe, is available on the Web that can be used to convert "banked" S-records to the linear S-record format required by the serial bootloader.

NOTE: *The bootloader is limited to receiving S-records containing a maximum of 64 bytes in the code/data field. If an S-record containing more than 64 bytes in the code/data field is received, an error message will be displayed.*

PPAGE Value	S-Record Address Range	Memory Type
\$00–\$2F	\$00000–\$BFFFF	Off-chip memory
\$30	\$C0000–\$C3FFF	On-chip FLASH
\$31	\$C4000–\$C7FFF	On-chip FLASH
\$32	\$C8000–\$CBFFF	On-chip FLASH
\$33	\$CC000–\$CFFFF	On-chip FLASH
\$34	\$D0000–\$D3FFF	On-chip FLASH
\$35	\$D4000–\$D7FFF	On-chip FLASH
\$36	\$D8000–\$DBFFF	On-chip FLASH
\$37	\$DC000–\$DFFFF	On-chip FLASH
\$38	\$E0000–\$E3FFF	On-chip FLASH
\$39	\$E4000–\$E7FFF	On-chip FLASH
\$3A	\$E8000–\$EBFFF	On-chip FLASH
\$3B	\$EC000–\$EFFFF	On-chip FLASH
\$3C	\$F0000–\$F3FFF	On-chip FLASH
\$3D	\$F4000–\$F7FFF	On-chip FLASH
\$3E	\$F8000–\$FBFFF	On-chip FLASH
\$3F	\$FC000–\$FFFFFF	On-chip FLASH

Figure 10. MC9S12DP256 PPAGE to S-Record Address Mapping

The conversion of the linear S-record load address to a PPAGE number and a PPAGE window address can be performed by the two formulas shown in [Figure 11](#). In the first formula, `PageNum` is the value written to the `PPAGE` register, `PPAGEWinSize` is the size of the PPAGE window which is \$4000. In the second formula, `PPAGEWinAddr` is the address within the PPAGE window where the S-record code/data is to be loaded. `PPAGEWinStart` is the beginning address of the PPAGE window which is \$8000.

```

pageNum = SRecLoadAddr / PPAGEWinSize;

PPAGEWinAddr = (SRecLoadAddr % PPAGEWinSize) + PPAGEWinStart;

```

Figure 11. PPAGE Number and Window Address Formulas

**Using
the S-Record
Bootloader**

The S-record bootloader presented in this application note utilizes the on-chip SCI for communications with a host computer and does not require any special programming software on the host.

The bootloader presented in this application note can be used to erase and reprogram all but the upper 4 K of on-chip FLASH memory. The bootloader program utilizes the on-chip SCI for communications and does not require any special programming software on the host computer. The only host software required is a simple terminal program that is capable of communicating at 9600 to 115,200 baud and supports XOn/XOff handshaking.

Invoking the bootloader causes the prompt shown in [Figure 12](#) to be displayed on the host terminal's screen. The lowercase ASCII characters a through c comprise the three valid bootloader commands. These three lowercase characters were selected, rather than the ASCII characters 1 through 3, to prevent accidental command execution. If a problem occurs while programming the FLASH, an error message is displayed, and the bootloader will redisplay its prompt and wait for a command entry from the operator. Because the host computer will continue sending the S-record file, each character of the S-record file would be interpreted as an operator command entry. Since S-records contain all of the ASCII numeric characters, it is highly likely that one of them would be understood as a valid command.

```

MC9S12DP256Bootloader

a) Erase Flash
b) Program Flash
c) Set Baud Rate
?

```

Figure 12. Serial Bootloader Prompt

Application Note

**Erase FLASH
Command**

Selecting the erase function by typing a lowercase `a` on the terminal will cause a bulk erase of all four 64-K FLASH arrays except for the 4-k boot block in the upper 64-K array where the S-record bootloader resides. After the erase operation is completed, a verify operation is performed to ensure that all locations were properly erased. If the erase operation is successful, the bootloader's prompt is redisplayed.

If any locations were found to contain a value other than `$FF`, an error message is displayed on the screen and the bootloader's prompt is redisplayed. If the MC9S12DP256 device will not erase after one or two attempts, the device may be damaged.

**Program FLASH
Command**

To increase the efficiency of the programming process, the S-record bootloader uses interrupt driven, buffered serial I/O in conjunction with XOn/XOff software handshaking to control the S-record data flow from the host computer. This allows the bootloader to continue receiving S-record data from the host computer while the data from the previously received S-record is programmed into the FLASH.

NOTE: *The terminal program must support XOn/XOff handshaking to properly reprogram the MC9S12DP256's FLASH memory.*

Typing a lowercase `b` on the terminal causes the bootloader to enter the programming mode, waiting for S-records to be sent from the host computer. The bootloader will continue to receive and process S-records until it receives an S8 or S9 end of file record. If the object file being sent to the bootloader does not contain an S8 or S9 record, the bootloader will not return its prompt and will continue to wait for the end of file record. Pressing the system's reset switch will cause the bootloader to return to its prompt.

If a FLASH memory location will not program properly, an error message is displayed on the terminal screen and the bootloader's prompt is redisplayed. If the MC9S12DP256 device will not program after one or two attempts, the device may be damaged or an S-record with a load address outside the range of the available on-chip FLASH may have been received. The S-record data must have load addresses in the range `$C0000-$FFFFFF`. This address range represents the upper 256 Kbytes of the 1-MB address space of the MC9S12DP256.

**Set Baud Rate
Command**

While the default communications rate of the bootloader is 9600 baud, this speed is much too slow if the majority of the MC9S12DP256's FLASH is to be programmed; however, it provides the best compatibility for initial communications with most terminal programs. The set baud rate command allows the bootloader communication rate to be set to one of four standard baud rates. Using a baud rate of 57,600 allows the entire 256 K of FLASH to be programmed in just under two minutes.

Typing a lowercase `c` on the terminal causes the prompt shown in **Figure 13** to be displayed on the host terminal's screen. Entering a number 1 through 4 on the keyboard will select the associated baud rate and issue a secondary prompt indicating that the terminal baud rate should be changed. After changing the terminal baud rate, pressing the enter or return key will return to the main bootloader prompt. The selected baud rate will remain set until the target system is reset.

```

1) 9600
2) 38400
3) 57600
4) 115200
? 3
Change Terminal BR, Press Return

```

Figure 13. Change Baud Rate Prompt

Bootloader Software

The software implementing the serial FLASH bootloader, shown in **Code Listing**, consists of seven basic parts: startup code, bootloader control loop, programming and erase code, serial communications routines, an S-record loader and a secondary interrupt vector jump table. The code is written in a position independent manner so that the generated object code will execute properly from any address.

Application Note

Startup Code

The bootloader startup code implements several setup and initialization tasks.

The first action performed by the startup code checks the state of pin 6 on port M. If a logic 1 is present, the `JMP` instruction will continue execution at the address stored in the reset vector of the secondary vector table. If a logic 0 is present at pin 6 of port M, execution continues at the label `Boot` where the COP watchdog timer is disabled.

After the watchdog timer is disabled, the bootloader copies itself into the upper 4 K of the on-chip RAM. Execution of the bootloader code from RAM is necessary so the portion of FLASH block zero not occupied by the bootloader can be erased and programmed. Notice that only the code between the labels `BootStart` and `BootLoadEnd` is copied into RAM. This does not include the secondary vector jump table or the primary interrupt vector addresses since neither is required by the bootloader. After the copy operation is complete, the RAM is relocated to overlay the upper 12 K of FLASH memory between `$D000` and `$FFFF`. Writes to the `INITRM` register do not go into effect until one bus clock after the write cycle occurs. This means that the RAM cannot be accessed at the new address until after this one clock delay. Normally, the store instruction would simply be followed with a NOP instruction to ensure that no unintended operations occurred. However, in this case because the RAM is being moved into the same address space where the CPU is executing, a CPU free cycle must follow the write cycle.

NOTE: *To understand why the store instruction must use extended addressing and must be aligned to an even byte boundary, it is necessary to examine the cycle-by-cycle execution detail of the store instruction.*

The `STAB` instruction using extended addressing requires three clock cycles when executed from internal MCU memory. These three clock cycles consist of a `P` cycle, a `w` cycle and an `O` cycle (`PwO`). The `P` cycle is a program word access cycle where program information is fetched as an aligned 16-bit word. The `w` cycle is the 8-bit data write. Finally, the `O` cycle is an optional cycle that is used to adjust instruction alignment in the instruction queue. An `O` cycle can be a free cycle (`£`) or a program word access cycle (`P`). When the first byte of an instruction with an odd number of bytes is misaligned (at an odd address), the `O` cycle becomes

a P cycle to maintain queue order. If the first byte is aligned (at an even address), the O cycle is an E cycle. Consequently, if the first byte of the STAB instruction using extended addressing is aligned to an even byte boundary, the O cycle will be an E cycle. This will then provide the cycle of delay required while the RAM is overlaying the FLASH. Because the default address of the INITRM register is in the direct page addressing range, most assemblers will use direct rather than extended addressing. The greater than character (>) appearing as the first character in the operand field of the STAB instruction is used to force extended addressing. Note that some assemblers may not recognize this modifier character.

The main reason for relocating the RAM, rather than executing the bootloader at the RAM's default address, is to allow the SCIO interrupt vector to be changed. Because the on-chip RAM has a higher priority in the memory decoding logic than the on-chip FLASH, overlaying the FLASH with the on-chip RAM causes the RAM to be accessed rather than the FLASH. Due to the fact that the bootloader's communications routines utilize the SCI in a buffered, interrupt driven mode, the SCIO interrupt vector must be initialized to point to the bootloader's SCI interrupt service routine.

After relocating the on-chip RAM, the startup code initializes the PLL and engages it as the bus clock. The values for the REFDV and SYNRR registers are calculated by the assembler based on values of the oscillator frequency (OscClk), final bus frequency (fEclock), and the desired reference frequency (RefClock). In this case, the final bus frequency is specified to be 24.0 MHz. Because this is an integer multiple of the oscillator frequency, the oscillator frequency can be used as the reference clock for the PLL. This results in a value of zero being written to the REFDV register. To obtain a bus clock of 24 MHz, the reference frequency must be multiplied by three. The value written to the SYNRR register multiplies the reference clock by SYNRR+1 to generate the bus clock. Therefore, a value of two is written to the SYNRR register to obtain a 24-MHz bus clock. Note that the four NOP instructions following the STAB instruction work around a bug in the 0K36N mask set. This errata manifested itself in the LOCK bit not being cleared until several bus cycles after a write to the SYNRR register had occurred. Also note that a 24-MHz bus clock was chosen to support a baud rate of 115,200.

The final actions performed by the startup code initialize the `FCLKDIV` register and call the `SCIInit` subroutine. The value written to the `FCLKDIV` register is calculated by the assembler and is based on the MC9S12DP256's oscillator frequency, not the bus frequency. The `SCIInit` subroutine initializes the SCI0 hardware and associated data structures needed to support buffered, interrupt driven communications. It accepts a single parameter in the D accumulator that is used to set the initial baud rate.

Bootloader Control Loop

After the startup code has completed its task, a sign-on message is displayed and the bootloader enters its main control loop. At the start of the loop, the X index register is loaded with the address of the bootloader prompt and the subroutine `PromptResp` is called. The `PromptResp` subroutine is used to display a null terminated (`$00`) character string and then waits for a single character response from the operator. Upon receipt of a character, the `PromptResp` subroutine returns and a range check is performed on the received character to ensure it is a valid command. If the received character is not a valid command, the entry is ignored and the prompt is redisplayed.

If the received character is one of the three valid commands, its ASCII value is used as an index into a table of offsets. However, before being used as an offset, the upper four bits of the ASCII value must be removed. Next, one must be subtracted from the remaining value because the first entry in the table is at an offset of zero. The result of the subtraction must then be multiplied by two because each entry in the table consists of two bytes. Next the `LEAX` instruction is used in conjunction with program counter relative (PCR) indexed addressing to load the address of the command table into the X index register in a position independent manner. Because the B accumulator contains an offset to the proper entry in the command table, the `LDD` instruction uses B accumulator offset indexed addressing to retrieve the entry from the table.

Examining the command table at label `CmdTable`, it can be seen that the table does not contain the absolute address of the command to execute. Rather each table entry contains an offset from the beginning of the table to the start of the command. This offset, when added to the

base address of the table contained in the X index register, produces the absolute address of the first instruction of the requested command. Using offsets in the command table in conjunction with calculating the beginning of the table in a position independent manner, allows a computed GOTO to be performed in a position independent manner. Finally, the JSR instruction uses accumulator offset indexed addressing to calculate the address of the command and calls the command as a subroutine.

Upon return from the command, the value of the global variable `ErrorFlag` is examined. If it contains a value of zero, the command completed without any errors. In this case, the code branches back to the top of the command loop where the bootloader prompt is redisplayed. If, however, an error occurred during command execution, the value in `ErrorFlag` is used as an index into a table of offsets to null terminated error strings. Calculation of the absolute address of the error string is performed in much the same manner as the calculation of the absolute address of the command. After displaying the error message, the code branches back to the top of the command loop where the bootloader prompt is redisplayed.

Program Command Code

The firmware required to implement the FLASH programming command consists of two subroutines. The first subroutine, `ProgFlash`, is called through the command table. This subroutine coordinates the activities required by the `ProgFBlock` subroutine which performs the actual programming of the FLASH memory. The `ProgFlash` subroutine begins by calling the `GetSRecord` subroutine which is used to receive a single S-record from the host computer. Having received an valid S-record, the subroutine performs several checks to ensure that the S-record meets the programming requirements of the MC9S12DP256. Because the MC9S12DP256's FLASH may only be programmed an align word at a time, both the code/data field length and the load address must be even numbers. If either value is odd, an error code is stored in the `ErrorFlag` global variable and the FLASH programming operation is terminated.

Next, the received S-record type is checked. Reception of an S8 or S9 S-record terminates the program FLASH command returning to the bootloader's control loop where the prompt is redisplayed. S0 records, designated as header records, do not contain any program or data and are simply ignored. Because the linear S-record addresses for the MC9S12DP256 begin at \$C0000 as shown in [Figure 10](#), only S2 S-records may be used to program the on-chip FLASH. Because the `GetSRecord` subroutine is capable of receiving S0, S1, S2, S8 and S9 S-records, the program FLASH command is terminated and an error code is returned in the `ErrorFlag` global variable if an S1 record is received.

After checking the received S-record type, a range check is performed on the S-record load address to ensure it is within the range of the on-chip FLASH minus the size of the 4 K protected area containing the bootloader. When performing the range check, the load address is first checked against `SRecLow`, the lowest valid S-record address for the on-chip FLASH. However, when checking against the upper limit, `SRecHi`, the number of code/data bytes contained in the S-record must be added to the load address before the comparison is performed. This ensures that even though the initial load address is less than the upper limit, none of the S-record code/data falls outside the upper limit.

Finally, the `ProgFlash` subroutine uses the S-record load address to calculate the PPAGE number and PPAGE window address using the formulas in [Figure 11](#). After initializing the PPAGE register, the PPAGE value is used to calculate a value for the block select bits. Closely examining the PPAGE values and the block numbers as shown in [Figure 1](#), it can be determined that the block number for any of the PPAGE values corresponds to the one's complement of bits two and three of the block's corresponding PPAGE value. After writing the proper value to the block select bits in the `FCNFG` register, the `ProgFBlock` subroutine is called to program the received S-record data into the FLASH. If no errors occurred during the programming operation, the code branches to the label `FSendPace` where an ASCII asterisk character is sent to the host computer to indicate that S-record data was successfully programmed into the FLASH.

The `ProgFBlock` subroutine performs the task of programming the received S-record data into the on-chip FLASH. While the subroutine generally follows the flowchart in [Figure 9](#), some operations have been rearranged to improve the efficiency of the implementation. The first two steps in the flowchart, writing the `PPAGE` register and block select bits, are performed in the `ProgFlash` subroutine. Note that the order of these two operations is not important. Because the value for the block select bits is derived from the `PPAGE` value, the `ProgFlash` subroutine writes the `PPAGE` register value first.

The third operation in the flowchart checks the state of the `CBEIF` bit to ensure that the command buffer is empty and ready to accept a new command. This check is not made at the beginning of the `ProgFBlock` subroutine because the bit is known to be set when the subroutine completes execution. This condition is inferred by the fact that the `CCIF` flag is set before the programmed data from the previously received S-record is verified.

The `ProgFBlock` subroutine begins by retrieving the S-record code/data field length, dividing the value by two and placing the result on the stack. The code/data field length is divided by two because the FLASH is programmed a word at a time. Next, the X and Y index registers are initialized to point to the FLASH and S-record data respectively. Note that the X index register is loaded with the value in the `PPAGEAddr` global variable. This value, calculated using the second formula in [Figure 11](#), will always point within the `PPAGE` window. After initializing the pointers, the programming loop is entered at label `ProgLoop`. Note that within the programming loop there are no instructions that directly correspond to the five bus cycle delay before checking the state of the `CBEIF` flag after issuing the program command. Instead, the five bus cycle delay is inherent in the three instructions (`LDAB`, `BITB`, `BNE`) used to check the state of the `ACCERR` and `PVIOL` status bits. This loop follows the remainder of the flowchart in [Figure 9](#), issuing a new programming command each time the `CBEIF` flag is set until all of the count in the local variable `NumWords` is zero.

Before verifying that all of the FLASH locations programmed properly, the firmware must wait until the `CCIF` flag is set, indicating that all issued programming commands have completed. Failure to observe this

constraint before performing a read operation on the FLASH will result in the setting of the ACCERR bit and any pending programming commands will be terminated. The verification process begins by reinitializing the `DataBytes` local variable and the X and Y index register pointers. If any of the programmed words do not match the S-record data, a “not equal” condition (Z bit in the CCR equal to 0) is returned.

Erase Command Code

The code comprising the FLASH erase command is not nearly as simple as the programming code; it consists of five subroutines. The reason for the additional complexity surrounds the method that must be used to erase a FLASH block containing protected areas. When a 64-K block has a portion of its contents protected from being erased or programmed, the FLASH’s mass erase command cannot be used. Instead, the unprotected areas must be erased one 512-byte sector at a time. Because the time required to erase a sector is 20 ms versus 100 ms for the mass erase operation, erasure of a 64-K block with protected areas requires much longer. In this case where the bootloader resides in a 4-K protected area of block zero, 120 sector erase operations must be performed. Not counting the time required to verify each sector erasure, the sector erase operations require 2.4 seconds (20 ms * 120 sectors).

The FLASH erase command begins with the subroutine `EraseFlash`, called through the command table. This subroutine coordinates the activities of the other four subroutines. It begins by performing a mass erase and verify on three of the 64-K FLASH blocks. After all three of the 64-K FLASH blocks have been successfully erased, the `EraseBlk0` subroutine is called to perform a sector by sector erase of the unprotected portion of FLASH block zero.

The `EraseBlk0` subroutine begins by allocating and initializing the local variable `PPAGECnt`. The initialized value of three is the number of 16-K PPAGE windows that will be completely erased a sector at a time. The PPAGE register is initialized with a value passed in the B accumulator from the `EraseFlash` subroutine. This value, \$3C, places the lower 16 K of FLASH block zero into the PPAGE window. The block select bits are initialized to zero. After loading the X index register with the address

of the start of the PPAGE window and the B accumulator with the number of sectors to erase, the `EraseSectors` subroutine is called. In addition to erasing the requested number of sectors, the `VerfSector` subroutine is called to verify the erasure. Note that the `VerfSector` subroutine verifies the erasure a word at a time because the erase verify command built into the FLASH state machine will only operate on a 64-K block. After `EraseBlk0` performs the erasure of the lower 48 K of FLASH block zero, the lower 24 sectors (\$8000–\$EFFF) of the upper 16 K of block zero are erased.

**Set Baud Rate
Command Code**

The code comprising the set baud rate command is relatively simple. The subroutine begins by displaying the baud rate change prompt and then waiting for the operator to enter a baud rate selection. A range check is performed on the entered character; if an invalid character is entered, the prompt is redisplayed. If the selection is valid, the upper four bits are masked off, one is subtracted from the lower four bits, and the result is divided by two. The result is used as an index into the `BaudTable` to retrieve the proper `SCI0BD` register value for the selected baud rate.

Before switching to the newly selected baud rate, a message is displayed prompting the operator to change the host terminal's baud rate. However, before the `SCI0BD` register is written with the new value, the firmware must wait until the last character of the message is shifted from the `SCI0` transmit shift register. Once the last character of the message is sent, the `SCI0BD` register is written with the new value and the `getchar` subroutine is called to wait for an indication from the operator that the host terminal baud rate has been changed. Finally, a carriage return/line feed is sent to the terminal before returning to the bootloader control loop.

**S-Record Loader
Code**

The `GetSRecord` subroutine is used to receive a single S-record from the host computer. `GetSRecord` begins by allocating space on the stack for two local variables and initializing the X index register. The `SRecBytes` variable is used to hold the converted value of the S-record length field. This value includes the number of bytes contained in the load address field, the length of the code/data field, and the length of the

checksum field. The variable `CheckSum` is used to contain the calculated checksum value as the S-record is received. The X index register is initialized to point to the beginning of the 24-bit global variable, `LoadAddr`, where the received S-record's address is stored. Note also that the most significant byte of `LoadAddr` is cleared in case an S1 record is received.

After the initializations, a search is begun for the character pairs S0, S1, S2, S8, or S9 which indicate the start of a valid S-record. Once a valid start of record is found, the number of bytes in the load address plus one is stored in the global variable `DataBytes`. This value is subsequently subtracted from the received S-record length byte to produce a result representing the code/data field length. Before receiving the S-record length byte, the second character of the start of record pair is stored in the global `RecType`. After receiving the S-record length byte, the value is saved in the local variable `SRecBytes`. This value is also used to initialize `CheckSum` which is used to calculate a checksum value as the S-record is received.

The loop beginning at the label `RcvData` receives the remainder of the S-record including the load address, the code/data field, and the checksum. Note that because each received byte is stored in successive memory locations, the global variables `LoadAddr` and `SRecData` must remain in the order they are declared. As each data byte and the checksum is received, it is added into the calculated checksum value. Because the received checksum is actually the one's complement of what the calculated checksum should be, adding the two values should produce a result of \$FF. incrementing the `CheckSum` variable at the end of the receive loop should produce a result of zero if the checksum and all the S-record fields were received properly. This results in an "equal" condition (CCR Z = 1) being returned if the S-record was properly received and a "not equal" condition (CCR Z = 0) being returned if a problem occurred receiving the S-record.

Operation of the `GetSRecord` subroutine is supported by the three additional subroutines `GetHexByte`, `IsHex`, and `CvtHex`. The `GetHexByte` subroutine retrieves two ASCII hex bytes from the serial port and converts them into a single 8-bit data byte that is returned in the B accumulator. The `IsHex` subroutine is used to check received byte to

ensure that it is an ASCII hexadecimal character. If the character in the B accumulator is a non-hexadecimal character, the subroutine returns a “not equal” condition (CCR Z = 0). Otherwise, an “equal” condition (CCR Z = 1) is returned. The `CvtHex` subroutine converts the ASCII hexadecimal character in the B accumulator to a binary value. The result remains in the B accumulator.

Serial Communications Code

The serial communications routines utilize SCI0 to communicate with a host computer. The routines utilize the SCI in an interrupt driven mode, allowing reception of data from the host computer while the bootloader is programming the on-chip FLASH memory. To prevent the possibility of the receive buffer overflowing, the receive routines support XOn/XOff handshaking with the host computer. Because the bootloader does not send large amounts of data to the host computer, XOn/XOff handshaking is not supported by the transmit routines.

To utilize the interrupt driven mode effectively, a circular buffer or queue must be associated with both the transmitter and receiver. The queue acts as an elastic buffer providing a software interface between the received character stream and the MC9S12DP256. In addition to the storage required by the transmit and receive queues, several other pieces of data are required for queue management. The information necessary to manage the queue consists of a way to determine the next available storage location in each queue, the next available location or piece of data in the queue, and a way to determine if a queue is full or empty. Rather than utilize 16-bit pointers to manage the queues, the communications routines employ four 1-byte variables. `RxIn`, `RxOut`, `TxIn`, and `TxOut` are used in conjunction with 8-bit accumulator offset indexed addressing to access data in the transmit and receive queues. In addition, two 1-byte variables, `RxBAvail` and `TxBAvail`, are used to keep track of the number of bytes available in each queue. When the value in each of these variables is equal to the size of the queue, the buffer is empty. When the value is zero, the queue is full. Using a byte for the index does not allow support of queue sizes greater than 255 bytes. However, this should not pose severe restrictions for most applications.

The proper queue size for an application will depend on the expected length of messages transmitted and received. If the selected transmit queue size is too small, the routines essentially will behave the same as the polled SCI example. Once the queue fills, the CPU12 will have to wait until a character is transmitted before the next character can be placed in the queue. If the receive queue is too small, there will be a risk that received characters will be lost if the queue becomes full and CPU12 does not remove some of the data before the next piece of data arrives. Conversely, picking queue sizes larger than necessary does not have a detrimental effect on program performance or loss of data. However, it will consume the valuable on-chip memory unnecessarily. If uncertain on the exact queue size for a particular application, it is best to make it larger than necessary. As shown, the transmit and receive queues do not have to be the same size, and their sizes are not required to be an even power of two.

The `XOffCount` and `XOnCount` constants are used to manage how full and how empty, respectively, the receive queue is allowed to get before the `XOff` and `XOn` control characters are sent to the host computer. The value for `XOffCount` should be chosen based on the number of bytes that are expected to be sent from the host after a request has been made for the `TxIRQ` routine to send an `XOff` to the host. This value, which represents the number of remaining bytes in the receive queue when an `XOff` should be sent, will depend on the UART characteristics of the host computer. In this case, a value of `XOffCount` would allow up to 10 additional characters to be sent after a request to send the `XOff` had been posted. This would allow for the host computer UART with an 8-byte FIFO plus the possible 2-character delay in sending the `XOFF` character if the transmit shift register and the transmit data register were both full.

The value for `XOnCount` should be selected such that the queue will never become empty as long as the host has data to send. Setting the correct value for this constant requires analysis of the rate at which data is removed from the queue by the application and the delay before the host computer begins sending data after receiving an `XOn`. Because the host's characteristics can vary widely, a value of the receive buffer minus eight was arbitrarily chosen. Note that the value of `XOnCount` represents the number of characters available in the receive queue.

The `SCIInit` subroutine is used to initialize the SCI hardware and the related queue data structures. The baud rate register (`SCI0BD`) value for the desired baud rate is passed to the subroutine in the D accumulator. The queue index values `RxIn`, `RxOut`, `TxIn`, `TxOut`, and the values for `RxBAvail` and `TxBAvail` are not specifically initialized by the subroutine because the initial values are set at the point of their declaration. This technique works in this case because the constant values were copied from the FLASH into RAM. In a situation where the variables were declared with a `ds` (define storage) directive each variable would have to be initialized to its proper value.

When the transmitter and receiver are enabled, notice that only the receive interrupts are enabled. Unlike the receiver interrupts, which may be enabled at all times, the transmit interrupt may be enabled only when the transmit queue contains characters to be sent. Enabling transmit interrupts at initialization would immediately cause a transmitter interrupt even though the transmit queue is empty. This is because the `TDRE` bit is set whenever the SCI transmitter is in an idle state. The final action performed by the `SCIInit` subroutine initializes the SCI0 interrupt vector to point to the SCI interrupt routine, `SCIISR`.

Because each SCI only has a single interrupt vector shared by the transmitter and receiver, a short dispatch routine determines the source of the interrupt and calls either the `RxIRQ` or `TxIRQ`. Note that it is not an arbitrary choice to have the dispatch routine check for receiver interrupts before transmitter interrupts. To avoid the loss of received data, an SCI interrupt dispatch routine should always check the receiver control and status flags before checking those associated with the transmitter. Failure to follow this convention will most likely result in receiver overruns when data is received during message transmissions longer than a couple of bytes.

The receive interrupt service routine, `RxIRQ`, has the responsibility of removing a received byte from the receive data register and placing it in the receive data queue if space is available. In addition, if space available in the queue falls below the value of `XOffCount`, two variables, `SendXOff` and `XOffSent`, are set to a non-zero value and transmitter interrupts are enabled. These actions cause an `XOff` character to be sent to the host computer the next time a transmit

interrupt is generated. `XOffSent` is used by the receive interrupt service routine to ensure that only a single `XOff` character is sent to the host after the space available in the queue falls below the value of `XOffCount`. `XOffSent` is also used by the `getchar` subroutine to determine if an `XOn` should be sent after each character is removed from the queue. Finally, notice that if the queue becomes full, the received byte is simply discarded.

The transmit interrupt service routine, `TxIRQ`, has the responsibility of removing a byte from the transmit data queue and sending it to the host computer. Before sending a character from the transmit queue, `SendXOff` is checked. If it contains a non-zero value, an `XOff` character is immediately sent to the host. Sending the `XOff` character before sending data that may be in the transmit queue ensures data flow from the host is stopped before the receive queue overflows. Notice that if the queue becomes empty after a character is transmitted, transmitter interrupts are disabled.

The last two major routines rounding out the serial communication code are the `getchar` and `putchar` subroutines. The `getchar` subroutine's main function is to retrieve a single character from the receive queue and return it to the calling routine in the B accumulator. Notice that if the receive queue is empty, the `getchar` subroutine will wait until a character is received from the host. Because this action may not be desirable for some applications, a utility subroutine, `SCIGetBuf`, can be called to determine if any data is in the receive queue. This small subroutine returns, in the B accumulator, a count of the number of data bytes in the receive queue. In addition to managing the receive queue variables each time a character is removed from the queue, the `getchar` subroutine checks the state of `XOffSent` and the number of characters left in the receive queue to determine if an `XOn` character should be sent to the host computer. If an `XOff` character was previously sent and the number of characters left in the receive queue is less than `XOnCount`, an `XOn` character is sent to the host by calling the `putchar` routine.

The `putchar` subroutine's main function is to place a single character, passed in the B accumulator, into the transmit queue. Once the character is in the queue and the queue variables have been updated, the transmit interrupt enable (TIE) bit is set. If transmitter interrupts were not previously enabled and the transmit data register empty (TDRE) bit is set, setting the TIE bit will cause an SCI interrupt to occur immediately.

Secondary Interrupt Vector Jump Table

Because the reset and interrupt vectors reside in the protected bootblock, a secondary vector table is located just below the protected bootblock area. Each entry in the secondary interrupt table should consist of a 2-byte address mirroring the primary interrupt and reset vector table. The secondary interrupt and reset vector table is utilized by having each vector point to a single JMP instruction that uses the CPU12's indexed-indirect program counter relative addressing mode. This form of the JMP instruction uses four bytes of memory and requires just six CPU clock cycles to execute. The table in [Figure 14](#) associates each vector source with the secondary interrupt table address.

Application Note

Interrupt Source	Secondary Vector Address	Interrupt Source	Secondary Vector Address
Reserved \$FF80	\$EF80	I ² C bus	\$EFC0
Reserved \$FF82	\$EF82	DLC	\$EFC2
Reserved \$FF84	\$EF84	SCME	\$EFC4
Reserved \$FF86	\$EF86	CRG lock	\$EFC6
Reserved \$FF88	\$EF88	Pulse accumulator B over o w	\$EFC8
Reserved \$FF8A	\$EF8A	Modulus down counter under o w	\$EFCA
PWM emergency shutdown	\$EF8C	Port H interrupt	\$EFCC
Port P interrupt	\$EF8E	Port J interrupt	\$EFCE
MSCAN 4 transmit	\$EF90	ATD1	\$EFD0
MSCAN 4 receive	\$EF92	ATD0	\$EFD2
MSCAN 4 errors	\$EF94	SCII	\$EFD4
MSCAN 4 wakeup	\$EF96	SCI0	\$EFD6
MSCAN 3 transmit	\$EF98	SPI0	\$EFD8
MSCAN 3 receive	\$EF9A	Pulse accumulator A input edge	\$EFDA
MSCAN 3 errors	\$EF9C	Pulse accumulator A over o w	\$EFD C
MSCAN 3 wakeup	\$EF9E	Timer over o w	\$EFD E
MSCAN 2 transmit	\$EFA0	Timer channel 7	\$EFE0
MSCAN 2 receive	\$EFA2	Timer channel 6	\$EFE2
MSCAN 2 errors	\$EFA4	Timer channel 5	\$EFE4
MSCAN 2 wakeup	\$EFA6	Timer channel 4	\$EFE6
MSCAN 1 transmit	\$EFA8	Timer channel 3	\$EFE8
MSCAN 1 receive	\$EFAA	Timer channel 2	\$EFEA
MSCAN 1 errors	\$EFAC	Timer channel 1	\$EFEC
MSCAN 1 wakeup	\$EFAE	Timer channel 0	\$EFEE
MSCAN 0 transmit	\$EFB0	Real-time interrupt	\$EFF0
MSCAN 0 receive	\$EFB2	IRQ	\$EFF2
MSCAN 0 errors	\$EFB4	XIRQ	\$EFF4
MSCAN 0 wakeup	\$EFB6	SWI	\$EFF6
FLASH	\$EFB8	Unimplemented instruction trap	\$EFF8
EEPROM	\$EFBA	COP failure reset	\$EFFA
SPI2	\$EFBC	Clock monitor fail reset	\$EFFC
SPI1	\$EFBE	Reset	\$EFFE

Figure 14. Secondary Vector Table Addresses for a 4-K Bootblock

Code Listing

```

00000000      equ      $0000
;
;
;
M      offset:
M      PCSave:
M      *
M      org      $:0
M      endm

;
M      switch:
M      ifc
M      ' .text', ':0'
M      org      PCSave
M      endif
M      endm

;
007A1200      OscClk:
016E3600      fEclock:
007A1200      RefClock:
00000000      REFVVal:
00000002      SYNVal:
00000000      if
00000000      equ      8000000
00000000      equ      24000000
00000000      equ      8000000
00000000      equ      (OscClk/RefClock)-1
00000000      equ      (fEclock/RefClock)-1
00000000      if
00000000      equ      (OscClk/200000/8)+FDIV8
00000000      else
00000028      FCLKDIVVal:
00000000      equ      (OscClk/200000)
endif

;
0000000D      Baud115200:
0000001A      Baud57600:
00000027      Baud38400:
0000009C      Baud9600:
;
00008000      FlashStart:
00001000      BootBlkSize:
00001000      RAMStart:
0000FF80      StackTop:
00003000      RAMBoot:
;
00000200      SectorSize:
00004000      PPAGESize:
;
000C0000      SRecLow:
000FF000      SRecHi:
;
; oscillator clock frequency.
; final E-clock frequency (PLL).
; reference clock used by the PLL.
; value for the REFV register.
; value for the SYN register.
; value for the FCLKDIV register.
; value for the FCLKDIV register.
; baud register value for 115,200 baud.
; baud register value for 57,600 baud.
; baud register value for 38,400 baud.
; baud register value for 9,600 baud.
; start address of the flash window.
; Erase protected bootblock size.
; default RAM base address.
; stack location after RAM is moved.
; starting RAM address where the bootloader
; will be copied.
; size of a Flash Sector.
; size of the PPAGE window ($8000 - $BFFF).
; lowest S-Record load address accepted
; by the bootloader.
; highest S-Record load address + 1
; accepted by the bootloader.

```

Application Note

```

00000030      S0RecType:      equ      '0'
00000031      S1RecType:      equ      '1'
00000032      S2RecType:      equ      '2'
00000038      S8RecType:      equ      '8'
00000039      S9RecType:      equ      '9'
;
00000001      FEraserError:    equ      1      ; Flash failed to erase.
00000002      SRecRngErr:      equ      2      ; S-Record out of range.
00000003      FlashPrgErr:    equ      3      ; Flash programming error.
00000004      SRecDataErr:    equ      4      ; Received S-Record contained an odd number
; of data bytes.
00000005      SRecAddrErr:    equ      5      ; S-Record Address is odd.
00000006      SRecLenErr:    equ      6      ; S-Record is too long.
;
;*****
;
;
;
;
0000F000      org      $f000

0000F000      BootStart:    brclr    PTIM,#$40,Boot      ; execute the bootloader?
0000F005      jmp          [Reset-BootBlkSize,pcr]      ; no. jump to the program pointed to by the
; secondary reset vector.
;
;
;
;
0000F009      Boot:        clr      COPCTL      ; keep watchdog disabled.
;
;
0000F00C      BootCopy:    lds      #StackTop      ; initialize the stack pointer
0000F00F      CEF000      ldx      #BootStart      ; point to the start of the Flash bootloader in Flash.
0000F012      CD3000      ldy      #RAMBoot      ; point to the start of on-chip RAM.
0000F015      CCF59A      ldd      #BootLoadEnd      ; calculate the size of the bootloader code.
0000F018      83F000      subd     #BootStart      ;
0000F01B      180A3070     movb   l,x+,l,y+      ; move a byte of the bootloader into RAM.
0000F01F      0434F9      dbne   d,MoveMore     ; dec byte count, move till done.
;
;
0000F022      C6C1      ldab   #$c0+RAMHAL     ; write to the INITRM register to overlay the Flash
; bootblock with RAM.
;
;
00000000      if      *%$0001<>0      ; PC currently at an odd byte boundary?
endif
;
;
0000F024      7B0010     stab   >INITRM        ; this instruction MUST use extended addressing an be
; aligned to an even byte boundary.
;
;
0000F027      C600      ldab   #REFDVVal      ; set the REFVDV register.
0000F029      5B35      stab   REFVDV         ;
0000F02B      C602      ldab   #SYNRVal      ; set the SYNR register.
0000F02D      5B34      stab   SYNR           ;
0000F02F      A7       nop                    ; nops required for bug in initial silicon.
0000F030      A7       nop
0000F031      A7       nop

```




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```

0000F032 A7      nop
0000F033 4F3708FC brclr CRGFLG,#LOCK,*
0000F037 4C3980    bset CLKSEL,#PLLSEL
;
0000F03A C628    ldab #FCLKDIVVal
0000F03C 7B0100    stab FCLKDIV
;
0000F03F CC009C    ldd #Baud9600
0000F042 15FA046F jsr SCIInit,pcr
0000F046 10EF      cli
;
0000F048 1AFA0055    leax SignOn,pcr
0000F04C 15FA0422    jsr OutStr,pcr
0000F050 69FA0546    clr ErrorFlag,pcr
0000F054 1AFA0064    leax BLPrompt,pcr
0000F058 072A    bsr PromptResp
0000F05A C161    cmpb #$61
0000F05C 25F2    blo CmdLoop
0000F05E C163    cmpb #$63
0000F060 22EE    bhi CmdLoop
0000F062 C40F    andb #$0f
0000F064 53    decb
0000F065 58    lslb
0000F066 1AFA0031    leax CmdTable,pcr
0000F06A ECE5    ldd b,x
0000F06C 15E6    jsr d,x
0000F06E E6FA0528 ldab ErrorFlag,pcr
0000F072 27DC    beq CmdLoop
0000F074 53    decb
0000F075 58    lslb
0000F076 1AFA00CC    leax ErrorTable,pcr
0000F07A ECE5    ldd b,x
0000F07C 1AE6    leax d,x
0000F07E 15FA03F0 jsr OutStr,pcr
0000F082 20CC    bra CmdLoop
;
;*****
;
PromptResp:
0000F084 15FA03EA    jsr OutStr,pcr
0000F088 15FA04B8    jsr getchar,pcr
0000F08C 15FA04E9    jsr putchar,pcr
0000F090 37    pshb
0000F091 1AFA0060    leax CrLfStr,pcr
0000F095 15FA03D9    jsr OutStr,pcr
0000F099 33    pulb
0000F09A 3D    rts
;
;*****
;

```



```

0000F21F 1AF903
0000F222 15FA024C
0000F226 4FCC40FC

0000F22A 3A
0000F22B 5CC8
0000F22D 15FA0313
0000F231 1AFAFEC0
0000F235 15FA0239
0000F239 3D

0000F23A 009C
0000F23C 0027
0000F23E 001A
0000F240 000D

;
BaudTable:
dc.w Baud9600
dc.w Baud38400
dc.w Baud57600
dc.w Baud115200

;
;*****
;
; ProgFlash:
equ *
ldab #PVIOL+ACCERR
stab FSTAT
bra FSkipFirst
ldab #'*'
jsr putchar,pcr
jsr GetSRecord,pcr
bne ProgDone
brclr DataBytes,pcr,#$01,DataLok ; if either the PVIOL or ACCERR bit is set from a
ldaa #SRECDataErr ; previous error, reset them so we can program the Flash.
bra ProgDone ; don't send the progress character the first time.
brclr LoadAddr+2,pcr,$$01,SReCOK ; the ascii asterisk is the progress character.
ldaa #SReCAddrErr ; let the user know we've processed an S-Record.
bra ProgDone ; go get an S-Record.
; non-zero condition means there was an error
brclr DataBytes,pcr,#$01,DataLok ; is the received S-Record length even?
ldaa #SRECDataErr ; no. report the error.
bra ProgDone ; stop programming.
brclr LoadAddr+2,pcr,$$01,SReCOK ; is the received S-Record address even?
ldaa #SReCAddrErr ; no. report the error.
bra ProgDone ; stop programming.

;
SReCOK:
ldab RecType,pcr ; check the record type.
cmpb #S1RecType ; S1 record received?
bne ChkNext ; no. check for S0, S2, S8 & S9 records.
ldaa #SReCngErr ; yes. only S2 records w/ load addresses $C0000 - $FEFFF
bra ProgDone ; allowed.
; save error & return.

;
ChkNext:
cmpb #S9RecType ; was it an S9 record?
beq ProgRtn ; yes. we're done.
cmpb #S8RecType ; was it an S8 record?
beq ProgRtn ; yes. we're done.
cmpb #S0RecType ; no. was it an S0 record?
beq FSendPace ; yes. just ignore it.

;
ldab LoadAddr,pcr ; was an S2 record. Get high byte of the 24-bit address.
cmpb #SReCLow>16 ; less than $c0000?
bhs ChkHiLimit ; no. check the upper limit.

```

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```

0000F289 8602 ; yes. S-Record out of range.
0000F28B 2045 ; save the error code & return.

0000F28D E6FA030B ; get the number of bytes in the S-Record.
0000F291 87 ; zero extend it.
0000F292 E3FA030A ; add in the lower 16-bits of the 24-bit address.
0000F296 B745 ; save rthe result in X.
0000F298 E6FA0303 ; get the upper 8-bits of the 24-bit address.
0000F29C C900 ; add in possible carry from lower 16-bits.
0000F29E C10F ; greater than $0fxxxx?
0000F2A0 2505 ; no. S-Record within range.
0000F2A2 8EF000 ; yes. check the lower 16- bits. Out of range?
0000F2A5 22E2 ; yes. S-Record out of range.

0000F2A7 E6FA02F4 ; get upper 8-bits of 24-bit load address.
0000F2AB B796 ; zero extend b into y for the 32-bit divide
0000F2AD ECFA02EF ; get the lower 16-bits of 24-bit load address.
0000F2B1 CE4000 ; divide the load address by the PPAGE window size.
0000F2B4 11
0000F2B5 C38000 ; add the PPAGE window start address to the remainder
; (this gives the PPAGE window load address).
; lower byte of the quotient is the PPAGE value.

0000F2B8 B7C6 ; calculate the value of the block select bits based
0000F2BA 5B30 ; on bits 3:2 of the PPAGE register value.
0000F2BC 54 ; mask off all but the lower 2 bits.
0000F2BD 54 ; select the block to erase.
0000F2BE 51 ; save the PPAGE window address.
0000F2BF C403 ; go program the data into Flash.
0000F2C1 7B0103 ; zero condition means all went ok.
0000F2C4 6DFA02D5 ; put error code where pod can access it.
0000F2C8 15FA000B ; if we fall through, we automatically return a non-zero condition.
0000F2CC 1827FF79 ;
0000F2D0 8603 ;
0000F2D2 6AFA02C4 ;
0000F2D6 3D ;

;*****
;
; offset 0
; set *
; org $0
;
; NumWords: ds 1
; LocalSize: set *
;
; switch .text
; ifc '.text', '.text'
; org PCsave
; endif

0000F2D7 E6FA02C1 ; ProgFLock:
; ldab DataBytes,pcr ; get the block size.

```



```

0000F2DB 54      lsrb          ; divide the byte count by 2 since we program a word
; at a time.
; allocate the local.
; get the PPAGE window Flash address.
; point to the received S-Record data.
; get a word from the buffer.
; latch the address & data into the Flash
; program/erase buffers.
; get the program command.
; write it to the command register.
; start the command by writing a 1 to CBEIF.
; check to see if there was a problem executing
; the command.
; if either the PVIOL or ACCERR bit is set,
; return.
; wait here till the command buffer is empty.
; any more words to program?
; yes. continue until done.
; no. wait until all commands complete.

0000F2DC 37      pshb
ldx      PPAGEWAddr,pcr
leay    SRecData,pcr
ldd     2,Y+
std     2,X+

0000F2E9 C620      ldab #PROG
stab    FCMD
ldab   #CBEIF
stab   FSTAT
ldab   FSTAT

0000F2F6 C530      bitb #PVIOL+ACCERR
bne    Return
brclr  FSTAT,#CBEIF,*
dec    NumWords,sp
bne    ProgLoop
brclr  FSTAT,#CCIF,*

;

0000F308 E6FA0290      ldab DataBytes,pcr
lsrb

0000F30D 6B80      stab NumWords,sp
ldx    PPAGEWAddr,pcr
leay   SRecData,pcr
ldd    2,Y+
cpd    2,X+
bne    Return
dec    NumWords,sp
bne    VerfLoop

;
Return: pulb
rts

;
;
;*****
;
offset 0
set *
org $0

;
BlockCnt: ds.b 1
;
LocalSize: set *
;
switch .text
ifc '.text','.text'
org PCSave

```

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Application Note

```

0000F323 C603
0000F325 37
0000F326 C630
0000F328 5B30

0000F32A 54
0000F32B 54
0000F32C 51
0000F32D C403
0000F32F 7B0103
0000F332 CE8000
0000F335 C641
0000F337 0760
0000F339 2621
0000F33B C605
0000F33D 075A
0000F33F 1F010540FB
0000F344 1E01050404
0000F349 8601
0000F34B 200F
0000F34D C604
0000F34F 7B0105
0000F352 D630
0000F354 CB04
0000F356 6380
0000F358 26CE
0000F35A 0706

0000F35C 6AFA023A
0000F360 33
0000F361 3D

0000F362
0000F362
00000000

00000000

00000001
00000001
00000001
0000F362

;
EraseFlash:      #S03
                  ltab
                  pshb
                  ldab
                  stab
                  #S30
                  PPAGE

EraseLoop:
                  lsrb
                  lsrb
                  comb
                  andb
                  #S03
                  FCNFG
                  stab
                  ldx
                  #FlashStart
                  ldab
                  #ERASE+WASS
                  bsr
                  EraseCmd
                  bne
                  SaveError
                  ldab
                  #ERVER+WASS
                  bsr
                  EraseCmd
                  brclr
                  FSTAT,#CCIF,*
                  brset
                  FSTAT,#BLANK,Erased
                  ldaa
                  #FEraseError
                  bra
                  SaveError
                  ldab
                  #BLANK
                  stab
                  FSTAT
                  ldab
                  PPAGE
                  addb
                  #S04
                  dec
                  BlockCnt,sp
                  bne
                  EraseLoop
                  bsr
                  EraseBlk0

SaveError:      staa
                 ErrorFlag,pcr
FEEDone:       pulb
                rts

;
;EraseBlk0 erases Flash block 0 a sector (512 bytes) at a time because the bootblock is protected.
;
offset 0
set *
org $0

PPAGECnt:     ds.b 1

;
LocalSize:    set *
;
switch .text
ifc '.text','.text'
org PCSave
endif
;

```

; write the PPAGE register to allow writes to the
; proper Flash block.
; calculate the value of the block select bits based
; on bits 3:2 of the PPAGE register.

; mask off all but the lower 2 bits.
; select the block to erase.
; latch address for erase command
; perform a bulk erase.

; if CCR Z=0, an error occurred.
; perform an erase verify.

; wait until the command has completed.
; flag a not erased error if the BLANK bit did not set.

; clear the BLANK status bit.

; get the current PPAGE value.
; add 4 to select the next 64K Flash block.
; done with 3 of the 64K blocks?
; no.

; block 0 must be erased seperately because it
; contains the bootblock.
; put error code where pod can access it.
; return.

; number of 16K PPAGE windows that will be
; completely erased.

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```

0000F362 1808AF03      EraseBlk0:      movb   #3,1,-sp
0000F366 5B30             stab   PPAGE
; 3 16K PPAGE windows will be completely erased.
; PPAGE for first 16K page of block 0
; (passed in the B accumulator).
; set block select bits to 0.
0000F368 790103      clr    FCNFG
0000F36B CE8000      EraseBlk0Loop: ldx   #FlashStart
0000F36E C620      #PPAGESize/SectorSize
0000F370 0712      bsr   EraseSectors
0000F372 260E      bne   BadBlk0
; point to the start of the PPAGE window.
; number of sectors in a PPAGE window.
; go erase the PPAGE window a sector at a time.
; non-zero value returned in A indicates a sector
; didn't erase.
; go to the next PPAGE.
0000F374 720030      inc   PPAGE
0000F377 6380      dec   PPAGECnt,sp
0000F379 26F0      bne   EraseBlk0Loop
0000F37B CE8000      ldx   #FlashStart
0000F37E C618      ldab  #((PPAGESize-BootBlkSize)/SectorSize ; number of sectors in PPAGE $3F
; minus the bootblock.
; erase all sectors outside the bootblock.
0000F380 0702      bsr   EraseSectors
0000F382 33             pulb  ; remove the page count from the stack.
0000F383 3D             rts
;
;Erases 'b' (accumulator) sectors beginning at address 'x' (index register)
;
EraseSectors:  exg   b,y
EraseSectLoop: ldab  #ERASE
; put the sector count in y.
; perform a sector erase.
bsr   EraseCmd
DoEraseVerf  beq   DoEraseVerf
; if no problem with the erase command, do a verify.
; if problem, return with an error code in a.
Rtn:         rts
DoEraseVerf:  bsr   VerfSector
; if problem, return with an error code in a.
bne   Rtn
SectorSize,x leax  SectorSize,x
dbne  Y,EraseSectLoop
; continue to erase remaining sectors.
; done. return.
rts
;
;Erases a block or sector of Flash
;
EraseCmd:     std   0,x
; latch address for erase command.
stab  FCMD
ldab  #CBEIF
stab  FSTAT
brclr FSTAT,#PVIOL+ACCERR,EraseCmdOK ; initiate the erase command.
; continue if the privilege violation &
; Access error flags are clear.
ldaa  #FEraseError
rts
EraseCmdOK:  brclr FSTAT,#CCIF,*
; wait until the command has completed.
clra
rts
;
;Verify that a sector was properly erased
;Must verify a word at a time because the built in verify command only works on a block (64K)
;

```

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Application Note

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```

0000F3B2 34      VerfSector:  pshx
0000F3B3 35      pshy
0000F3B4 CD0100      ldy
0000F3B7 EC31      #SectorSize/2
0000F3B9 048404      ldd
0000F3BC 8601      d,WordOK
0000F3BE 2004      #FFraseError
0000F3C0 0436F4      bra SectRtn
0000F3C3 87      WordOK:    dbne Y,VerfSectLoop
0000F3C4 31      clra
0000F3C5 30      SectRtn:  pulx
0000F3C6 3D      rts
;
;*****
;
; offset 0
; set *
; org $0
;
; SRecBytes: ds.b 1
; CheckSum:  ds.b 1
; LocalSize: set *
;
; switch .text
; ifc '.text','.text'
; org PCSave
; endif
;
; GetSRecord:
; equ *
; leas -LocalSize,sp
; leax LoadAddr,pcr
; clr 0,x
;
; LookForSOR:
; jsr getchar,pcr
; cmpb #'S'
; bne LookForSOR
; jsr getchar,pcr
; cmpb #S0Rectype
; bne CheckForS9
; bra Addr16
;
; CheckForS9:
; cmpb #S9Rectype
; bne ChkForS1
; bra Addr16
;
; ChkForS1:
; cmpb #S1Rectype
; bne ChkForS2
;
; Addr16:   inx
; save the base address of the sector.
; save the sector count.
; we'll check 2 bytes at a time.
; get a byte from the sector.
;
; yes. dec the sector word count.
; restore the sector count.
; restore the base address of the sector.
; return.
; number of bytes in the address, data & checksum fields.
; used for calculated checksum.
; allocate stack space for variables.
; point to the code/data buffer.
; clear the upper byte of the 24 bit address
; (in case we receive a 16-bit address).
; get a character from the receiver.
; start-of-record character?
; no. go back & get another character.
; yes. we found the start-of-record character (ASCII 'S')
; found an S0 (header) record?
; no. go check for an S9 record.
; yes. go receive the S0 record. (16-bit load address)
; found an S9 (end) record? (16-bit load address)
; no. go check for an S1 record.
; go receive the S9 record.
; found an S1 record? (16-bit load address)
; no. false start-of-record character received.
; go check for another.
; adjust the storage pointer to compensate for

```



```

0000F3EC 8603          ldaa      #3
0000F3EE 6AFA01AA        staa     DataBytes,pcr
0000F3F2 2010          bra      SaveRecType
;
0000F3F4 C132          cmpb     #S2RecType
0000F3F6 2602          bne     ChkForS8
0000F3F8 2004          bra     Addr24
;
0000F3FA C138          cmpb     #S8RecType
0000F3FC 26D1          bne     LookForSOR
0000F3FE 8604          ldaa     #4
0000F400 6AFA0198        staa     DataBytes,pcr
0000F404 6BFA0193        stab    RecType,pcr
;
0000F408 15FA003E        jsr     GetHexByte,pcr
0000F40C 2626          bne     BadSRec
0000F40E 6B80          stab    SRecBytes,sp
;
0000F410 6B81          stab    CheckSum,sp
;
0000F412 E0FA0186        subb    DataBytes,pcr
;
0000F416 6BFA0182        stab    DataBytes,pcr
0000F41A C140          cmpb     #64
0000F41C 2304          bls     RcvData
0000F41E 8606          ldaa     #SRecLenErr
0000F420 2012          bra     BadSRec
0000F422 15FA0024        jsr     GetHexByte,pcr
0000F426 260C          bne     BadSRec
0000F428 6B30          stab    1,x+
0000F42A EB81          addb    CheckSum,sp
0000F42C 6B81          stab    CheckSum,sp
0000F42E 6380          dec     SRecBytes,sp
0000F430 26F0          bne     RcvData
0000F432 6281          inc     CheckSum,sp
0000F434 1B82          leas   LocalSize,sp
0000F436 3D          rts
;
;*****
;
IsHex:
0000F437          equ     *
0000F437 C130          cmpb    #'0'
0000F439 250E          blo     NotHex
;
0000F43B C139          cmpb    #'9'
0000F43D 2308          bls     ISHex1
0000F43F C141          cmpb    #'A'
0000F441 2506          blo     NotHex
;
; less than ascii hex zero?
; yes. character is not hex. return a non-zero
; ccr indication.
; less than or equal to ascii hex nine?
; yes. character is hex. return a zero ccr indication.
; less than ascii hex 'A'?
; yes. character is not hex. return a non-zero
; ccr indication.
;*****

```

Application Note

```

0000F443 C146      ; less than or equal to ascii hex 'F'?
0000F445 2202      ; yes. character is hex. return a non-zero
                   ; ccr indication.
0000F447 1404      ; no. return a zero ccr indication.
0000F449 3D

IsHex1:
NotHex:
;
;
;*****
;
;
GetHexByte:
equ      *
jsr     getchar,pcr
bsr     IsHex
beq     OK1
rts
bsr     CvtHex
ldaa   #16
mul
pshb
jsr     getchar,pcr
bsr     IsHex
beq     OK2
pulb
rts
bsr     CvtHex
addb   1,sp+
clra
rts
OK1:
OK2:
;
;
;*****
;
;
CvtHex:
subb   #'0'
cmpb   #09
bls    CvtHexRtn
subb   #07
rts
CvtHexRtn:
;
;
;*****
;
;
OutStr:
equ      *
ldab   1,x+
beq     OutStrDone
jsr     putchar,pcr
bra    OutStr
rts
OutStrDone:
;
;
;*****
;
;
RxBufSize: equ 32
TxBufSize: equ 16
;

```



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```
00000018 RxOnCount: equ RxBufSize-8
0000000A XOffCount: equ 10
00000011 XOn: equ $11
00000013 XOff: equ $13
0000F47D RxBuf: dcb RxBufSize,0
0000F49D TxBuf: dcb TxBufSize,0
0000F4AD RxIn: dcb 0
0000F4AE RxOut: dcb 0
0000F4AF TxIn: dcb 0
0000F4B0 TxOut: dcb 0
0000F4B1 RxBAvail: dcb RxBufSize
0000F4B2 TxBAvail: dcb TxBufSize
0000F4B3 XOffSent: dcb 0
0000F4B4 SendXOff: dcb 0
;
;
;*****
; SCIInit: std SCI0BD
; ldab #TE+RE+RIE
; stab SCI0CR2
; leax SCIISR_pcr
; stx SCI0
; rts
;*****
; SCIISR: brclr SCI0CR2,#RIE,ChkRxInts
; brset SCI0SR1,#RDRF,RxIRQ
; brclr SCI0CR2,#TIE,NoSCIInt
; brset SCI0SR1,#TDRE,TxIRQ
; rti
;*****
; RxIRQ: tst XOffSent,pcr
; bne AlreadySent
; ldaa RxBAvail,pcr
; cmpa #XOffCount
; bhi AlreadySent
; inc SendXOff
; bset SCI0CR2,#TIE
; inc XOffSent,pcr
;
; number of bytes avail. in the Rx queue
; before an XOn can be sent.
; number of bytes remaining in the Rx queue
; when an XOff is sent.
; ASCII DC1
; ASCII DC3
; receive queue.
; transmit queue.
; next available location in the Rx queue.
; next character to be removed from the Rx queue.
; next available location in the Tx queue
; next character to be sent from the Tx queue.
; number of bytes left in the Rx queue.
; number of bytes left in the Tx queue.
; if != 0, an XOff has been sent.
; request to TX ISR to send an XOff to the host if != 0.
```

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```

0000F543 3D      rts
;
;*****
;
; return number available.
;
; save the registers we'll use.
;
; any characters available?
; no. just wait until some are.
; point to the physical start of the Rx queue.
; get the index to the next available character
; in the Rx queue.
; get the character.
; point to the next location in the queue.
; reached the end of the queue?
; no.
; yes. wrap to the start.
; update the queue index.
; we removed a character from the queue, there's
; 1 more available.
; was an XOff character previously sent by the RX ISR?
; no. just return.
; yes. get the number of bytes available in the Rx queue.
; enough space available to receive more?
; no. just return.
; yes. save the character we retrieved from the Rx queue.
; send an XOn character to the host.
; clear the XOff flag.
; restore the character we retrieved from the Rx queue.
; restore what we saved.
;
; return.
;
;*****
;
; save the registers we'll use.
; Any room left in the Tx queue?
; no. just wait here till it is.
; point to the physical start of the Tx queue.
; get the index to the next available spot.
; put the character in.
; point to the next available spot.
; go past the end of the queue?
; no.
; yes. wrap around to the start.
; update the queue index
; one less byte available in the Tx queue.
; enable transmitter interrupts.
;
;*****
;
;*****
;
getchar:      pshx
psha
ldab
RxChk:      #RxBufSize
             RxBAvail,pcr
             RxChk
             RxBuff,pcr
             RxOut,pcr
             a,x
             inca
             cmpa
             blo
             clra
             staa
             inc
             NopcWrap:
             tst
             beq
             ldaa
             cmpa
             bhs
             pshb
             ldab
             bsr
             clr
             pulb
             pula
             pulx
             rts
;
;*****
;
putchar:      pshx
psha
TxChk:      TxBAvail,pcr
             TxChk
             leax
             ldaa
             stab
             inca
             cmpa
             blo
             clra
             staa
             dec
             bset
             NopcWrap:
             TxIn,pcr
             TxBAvail,pcr
             SCT0CR2,#TIE
;
;*****
;
;*****
;

```

Application Note

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```

0000F597 32      pula
0000F598 30      pulx
0000F599 3D      rts

;
;
; restore what we saved.
;
; return

BootLoadEnd: equ *

;
;
; Global Variable declarations
;
;
; ErrorFlag:      ds.b 1
; RecType:       ds.b 1
;
; DataBytes:     ds.b 1
; PPAGEWAddr:   ds.b 2
; LoadAddr:     ds.b 3
; SRecData:     ds.b 65

;
;
; *****
; This is the jump table that is used to access the secondary interrupt vector table. Each one
; of the actual interrupt vectors, beginning at $ff8c, points to an entry in this table. Each jmp
; instruction uses indexed indirect program counter relative (pcr) addressing to access the
; secondary interrupt vector table that is located just below the bootblock.
;
; *****
;
JPWMEShutdown: jmp [PWMEShutdown-BootBlkSize,pcr]
JPortPInt:     jmp [PortPInt-BootBlkSize,pcr]
JMSCAN4Tx:     jmp [MSCAN4Tx-BootBlkSize,pcr]
JMSCAN4Rx:     jmp [MSCAN4Rx-BootBlkSize,pcr]
JMSCAN4Errs:   jmp [MSCAN4Errs-BootBlkSize,pcr]
JMSCAN4WakeUp: jmp [MSCAN4WakeUp-BootBlkSize,pcr]
JMSCAN3Tx:     jmp [MSCAN3Tx-BootBlkSize,pcr]
JMSCAN3Rx:     jmp [MSCAN3Rx-BootBlkSize,pcr]
JMSCAN3Errs:   jmp [MSCAN3Errs-BootBlkSize,pcr]
JMSCAN3WakeUp: jmp [MSCAN3WakeUp-BootBlkSize,pcr]
JMSCAN2Tx:     jmp [MSCAN2Tx-BootBlkSize,pcr]
JMSCAN2Rx:     jmp [MSCAN2Rx-BootBlkSize,pcr]
JMSCAN2Errs:   jmp [MSCAN2Errs-BootBlkSize,pcr]
JMSCAN2WakeUp: jmp [MSCAN2WakeUp-BootBlkSize,pcr]
JMSCAN1Tx:     jmp [MSCAN1Tx-BootBlkSize,pcr]
JMSCAN1Rx:     jmp [MSCAN1Rx-BootBlkSize,pcr]
JMSCAN1Errs:   jmp [MSCAN1Errs-BootBlkSize,pcr]
JMSCAN1WakeUp: jmp [MSCAN1WakeUp-BootBlkSize,pcr]
JMSCAN0Tx:     jmp [MSCAN0Tx-BootBlkSize,pcr]

```




```

0000F62F 05FBF97F JMSCAN0Rx: jmp [MSCAN0Rx-BootBlkSize,pcr]
0000F633 05FBF97D JMSCAN0Errs: jmp [MSCAN0Errs-BootBlkSize,pcr]
0000F637 05FBF97B JMSCAN0WakeUp: jmp [MSCAN0WakeUp-BootBlkSize,pcr]
0000F63B 05FBF979 JFlash: jmp [Flash-BootBlkSize,pcr]
0000F63F 05FBF977 JEEPROM: jmp [EEPROM-BootBlkSize,pcr]
0000F643 05FBF975 JSP12: jmp [SP12-BootBlkSize,pcr]
0000F647 05FBF973 JSP11: jmp [SP11-BootBlkSize,pcr]
0000F64B 05FBF971 JIICBus: jmp [IICBus-BootBlkSize,pcr]
0000F64F 05FBF96F JDLC: jmp [DLC-BootBlkSize,pcr]
0000F653 05FBF96D JSCME: jmp [SCMEVect-BootBlkSize,pcr]
0000F657 05FBF96B JCRGLock: jmp [CRGLock-BootBlkSize,pcr]
0000F65B 05FBF969 JPACCBov: jmp [PACCBov-BootBlkSize,pcr]
0000F65F 05FBF967 JModDnCtr: jmp [ModDnCtr-BootBlkSize,pcr]
0000F663 05FBF965 JPortHint: jmp [PortHint-BootBlkSize,pcr]
0000F667 05FBF963 JPortJInt: jmp [PortJInt-BootBlkSize,pcr]
0000F66B 05FBF961 JATD1: jmp [ATD1-BootBlkSize,pcr]
0000F66F 05FBF95F JATD0: jmp [ATD0-BootBlkSize,pcr]
0000F673 05FBF95D JSC11: jmp [SC11-BootBlkSize,pcr]
0000F677 05FBF95B JSC10: jmp [SC10-BootBlkSize,pcr]
0000F67B 05FBF959 JSP10: jmp [SP10-BootBlkSize,pcr]
0000F67F 05FBF957 JPACCAEdge: jmp [PACCAEdge-BootBlkSize,pcr]
0000F683 05FBF955 JPACCAov: jmp [PACCAov-BootBlkSize,pcr]
0000F687 05FBF953 JTimerOv: jmp [TimerOv-BootBlkSize,pcr]
0000F68B 05FBF951 JTimerCh7: jmp [TimerCh7-BootBlkSize,pcr]
0000F68F 05FBF94F JTimerCh6: jmp [TimerCh6-BootBlkSize,pcr]
0000F693 05FBF94D JTimerCh5: jmp [TimerCh5-BootBlkSize,pcr]
0000F697 05FBF94B JTimerCh4: jmp [TimerCh4-BootBlkSize,pcr]
0000F69B 05FBF949 JTimerCh3: jmp [TimerCh3-BootBlkSize,pcr]
0000F69F 05FBF947 JTimerCh2: jmp [TimerCh2-BootBlkSize,pcr]
0000F6A3 05FBF945 JTimerCh1: jmp [TimerCh1-BootBlkSize,pcr]
0000F6A7 05FBF943 JTimerCh0: jmp [TimerCh0-BootBlkSize,pcr]
0000F6AB 05FBF941 JRTI: jmp [RTI-BootBlkSize,pcr]
0000F6AF 05FBF93F JIRQ: jmp [IRQ-BootBlkSize,pcr]
0000F6B3 05FBF93D JXIRQ: jmp [XIRQ-BootBlkSize,pcr]
0000F6B7 05FBF93B JSWI: jmp [SWI-BootBlkSize,pcr]
0000F6BB 05FBF939 JIllop: jmp [Illop-BootBlkSize,pcr]
0000F6BF 05FBF937 JCOPFail: jmp [COPFail-BootBlkSize,pcr]
0000F6C3 05FBF935 JClockFail: jmp [ClockFail-BootBlkSize,pcr]
;
0000FF0D org $ff0d
;
0000FF0D CF dc.b ; setup a 4K bootblock in Flash block 0.
;
0000FF0F org $ff0f ; location of security byte.
;
0000FF0F FE dc.b ; value of security byte for unsecured state.
;
0000FF8C org $ff8c
;
0000FF8C F5E3 dc.w JPWMEShutdown

```

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0000FF8E F5E7	PortPInt:	dc.w	JPortPInt
0000FF90 F5EB	MSCAN4Tx:	dc.w	MSCAN4Tx
0000FF92 F5EF	MSCAN4Rx:	dc.w	MSCAN4Rx
0000FF94 F5F3	MSCAN4Errs:	dc.w	MSCAN4Errs
0000FF96 F5F7	MSCAN4WakeUp:	dc.w	MSCAN4WakeUp
0000FF98 F5FB	MSCAN3Tx:	dc.w	MSCAN3Tx
0000FF9A F5FF	MSCAN3Rx:	dc.w	MSCAN3Rx
0000FF9C F603	MSCAN3Errs:	dc.w	MSCAN3Errs
0000FF9E F607	MSCAN3WakeUp:	dc.w	MSCAN3WakeUp
0000FFA0 F60B	MSCAN2Tx:	dc.w	MSCAN2Tx
0000FFA2 F60F	MSCAN2Rx:	dc.w	MSCAN2Rx
0000FFA4 F613	MSCAN2Errs:	dc.w	MSCAN2Errs
0000FFA6 F617	MSCAN2WakeUp:	dc.w	MSCAN2WakeUp
0000FFA8 F61B	MSCAN1Tx:	dc.w	MSCAN1Tx
0000FFAA F61F	MSCAN1Rx:	dc.w	MSCAN1Rx
0000FFAC F623	MSCAN1Errs:	dc.w	MSCAN1Errs
0000FFAE F627	MSCAN1WakeUp:	dc.w	MSCAN1WakeUp
0000FFB0 F62B	MSCAN0Tx:	dc.w	MSCAN0Tx
0000FFB2 F62F	MSCAN0Rx:	dc.w	MSCAN0Rx
0000FFB4 F633	MSCAN0Errs:	dc.w	MSCAN0Errs
0000FFB6 F637	MSCAN0WakeUp:	dc.w	MSCAN0WakeUp
0000FFB8 F63B	Flash:	dc.w	JFlash
0000FFBA F63F	EEPROM:	dc.w	JEEPROM
0000FFBC F643	SPI2:	dc.w	JSPI2
0000FFBE F647	SPI1:	dc.w	JSPI1
0000FFC0 F64B	IICBus:	dc.w	JIICBus
0000FFC2 F64F	DLC:	dc.w	JDLC
0000FFC4 F653	SCMEVect:	dc.w	JSCME
0000FFC6 F657	JCRGLock:	dc.w	JCRGLock
0000FFC8 F65B	PACCBov:	dc.w	JPACCBov
0000FFCA F65F	ModDnCtr:	dc.w	JModDnCtr
0000FFCC F663	PortHInt:	dc.w	JPortHInt
0000FFCE F667	PortJInt:	dc.w	JPortJInt
0000FFD0 F66B	ATD1:	dc.w	JATD1
0000FFD2 F66F	ATD0:	dc.w	JATD0
0000FFD4 F673	SCI1:	dc.w	JSCI1
0000FFD6 F677	SCI0:	dc.w	JSCI0
0000FFD8 F67B	SPI0:	dc.w	JSPI0
0000FFDA F67F	PACCAEdge:	dc.w	JPACCAEdge
0000FFDC F683	PACCAov:	dc.w	JPACCAov
0000FFDE F687	TimerOv:	dc.w	JTimerOv
0000FFE0 F68B	TimerCh7:	dc.w	JTimerCh7
0000FFE2 F68F	TimerCh6:	dc.w	JTimerCh6
0000FFE4 F693	TimerCh5:	dc.w	JTimerCh5
0000FFE6 F697	TimerCh4:	dc.w	JTimerCh4
0000FFE8 F69B	TimerCh3:	dc.w	JTimerCh3
0000FFEA F69F	TimerCh2:	dc.w	JTimerCh2
0000FFEC F6A3	TimerCh1:	dc.w	JTimerCh1
0000FFEE F6A7	TimerCh0:	dc.w	JTimerCh0
0000FFF0 F6AB	RTI:	dc.w	JRTI



```

0000FFF2 F6AF      IRQ:      JIRQ
0000FFF4 F6B3     XIRQ:     JXIRQ
0000FFF6 F6B7     SWI:      JSWI
0000FFF8 F6BB     I11op:    JI11op
0000FFFA F6BF     COPFail:  JCOPFail
0000FFFC F6C3     ClcokFail: JClcokFail
0000FFFE F000     Reset:    dc.w  BootStart

Errors: None
Labels: 472
Last Program Address: $0000FFFF
Last Storage Address: $FFFFFFFF
Program Bytes: $000006F4 1780
Storage Bytes: $0000004E 78

```

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